

Nullstellensatz Size-Degree Trade-offs from Reversible Pebbling

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University of Copenhagen and KTH Royal Institute of Technology

AC Section lunch meeting

December 17, 2019



Joint work with Susanna F. de Rezende, Or Meir, and Robert Robere

Nullstellensatz (NS)

- ▶ Polynomials $\{P_1 = 0, P_2 = 0, \dots, P_m = 0\}$ in $\mathbb{F}[x_1, \dots, x_n]$

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- ▶ Degree: maximum degree (3 in example)
- ▶ Size: # monomials when expanded (7 in example)

Questions of interest

- ▶ Upper and lower bounds on degree

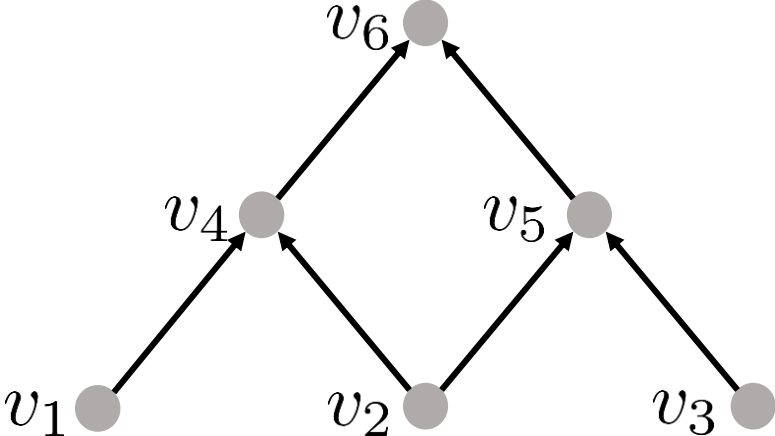
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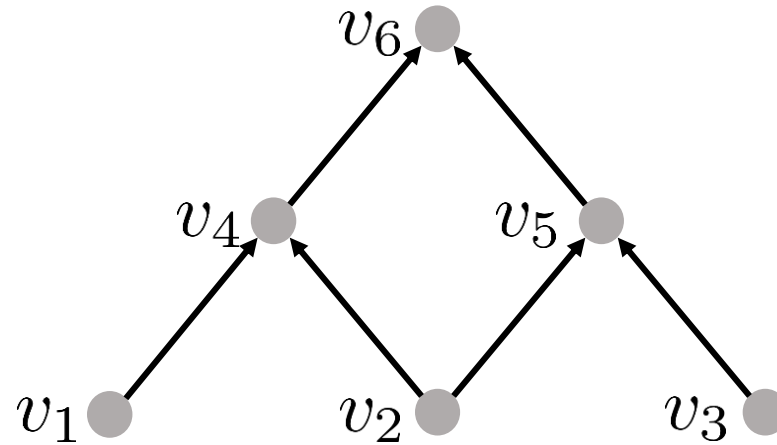
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- ▶ Size-degree relations? Simultaneous optimization?

Reversible pebble game



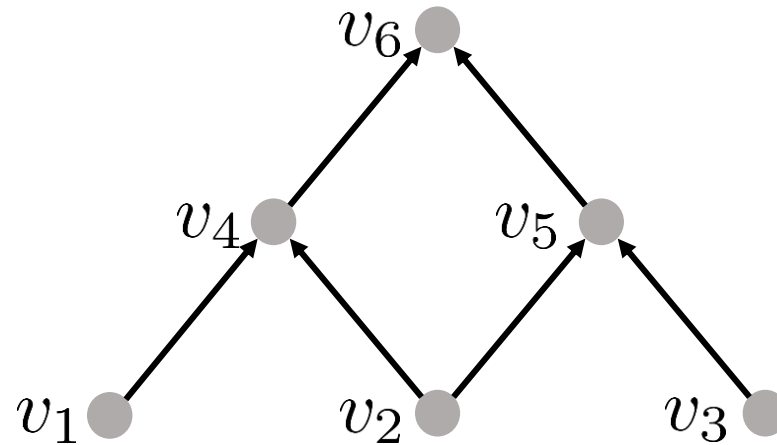
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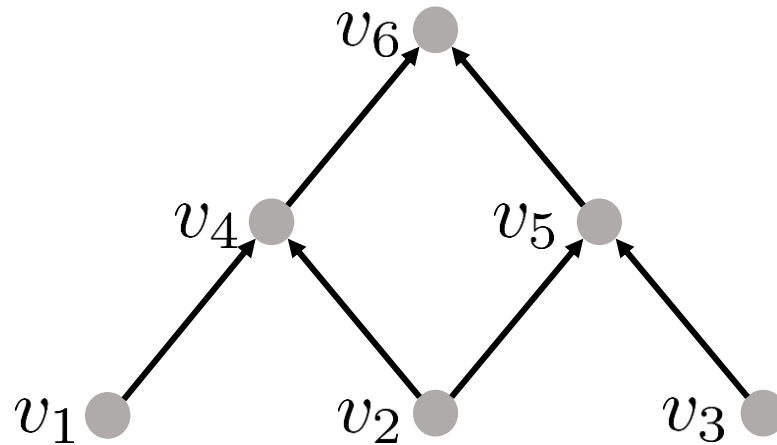
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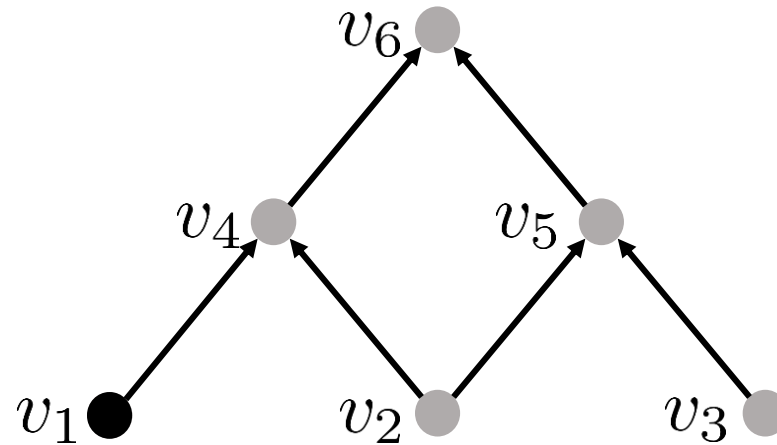
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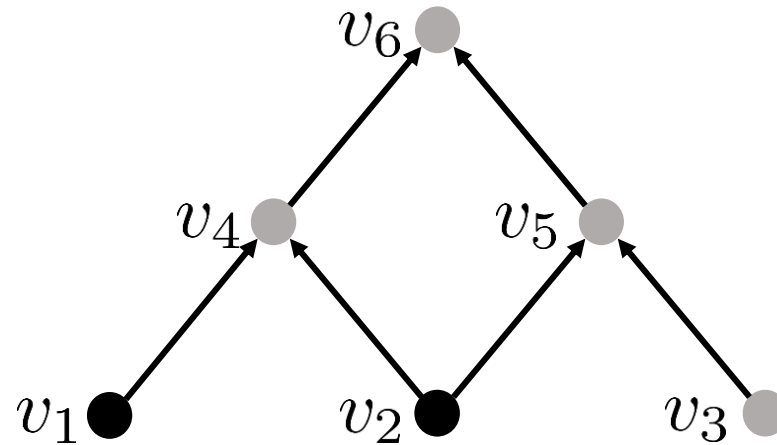


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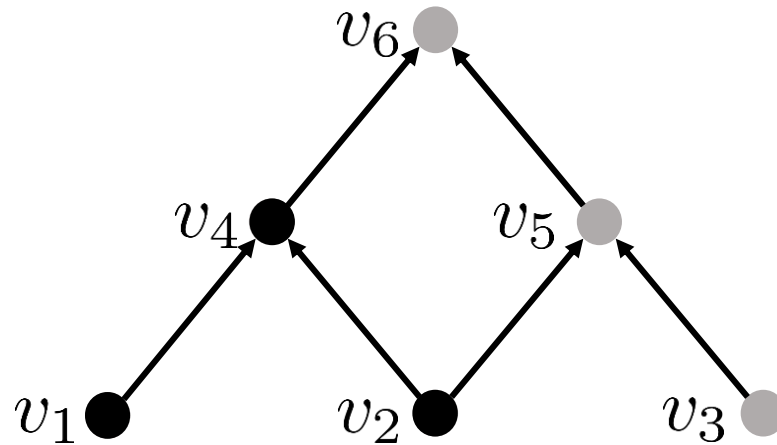


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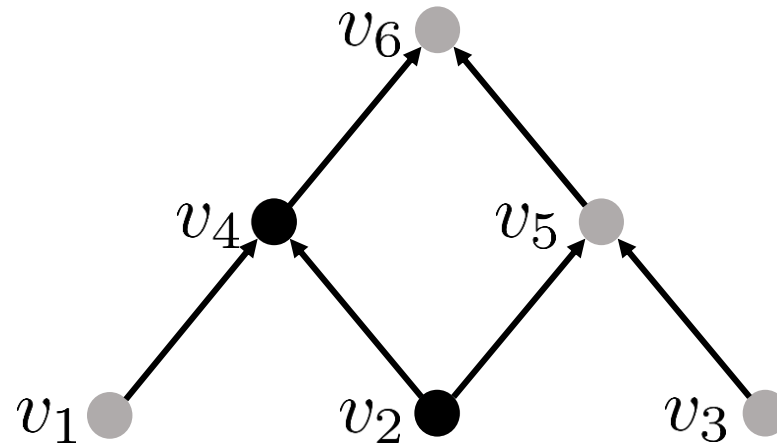
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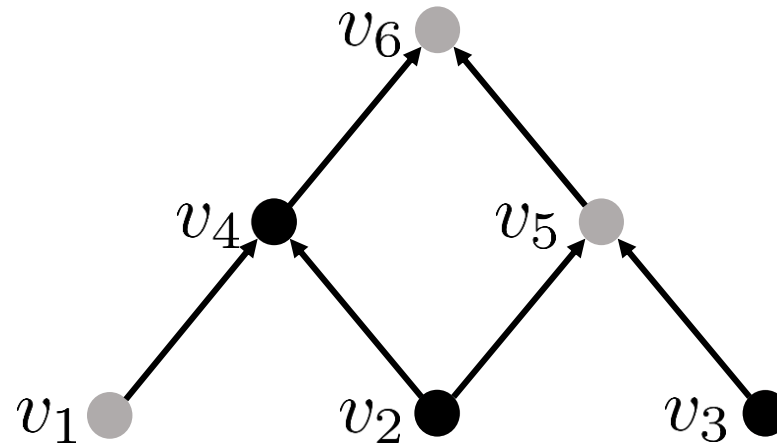


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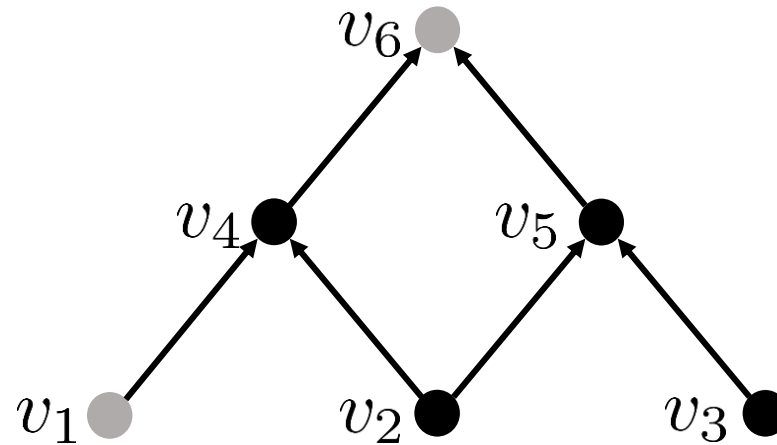


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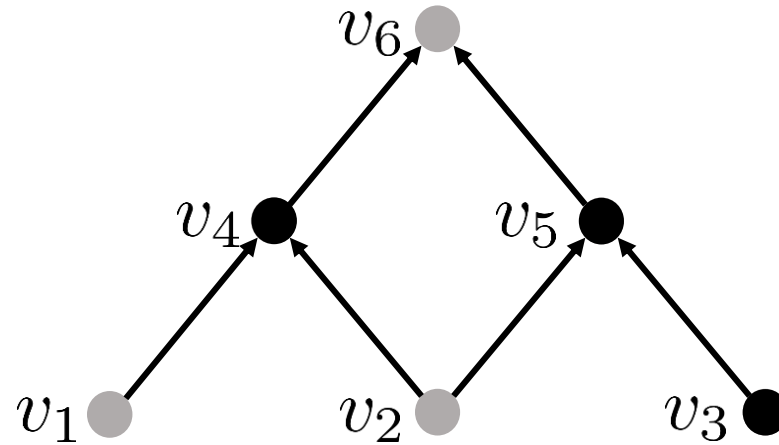
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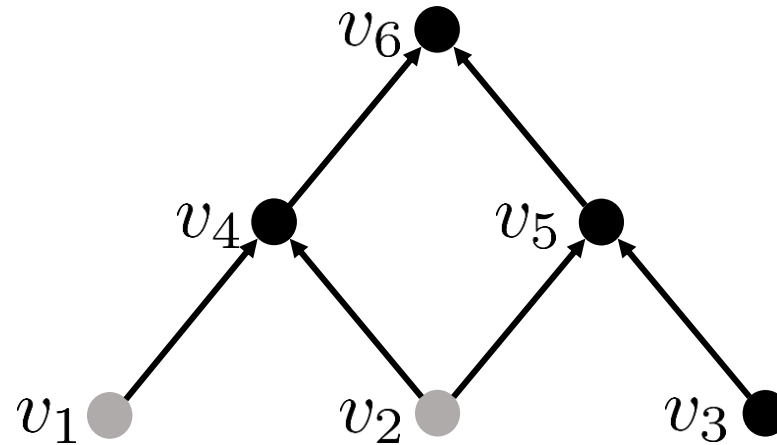
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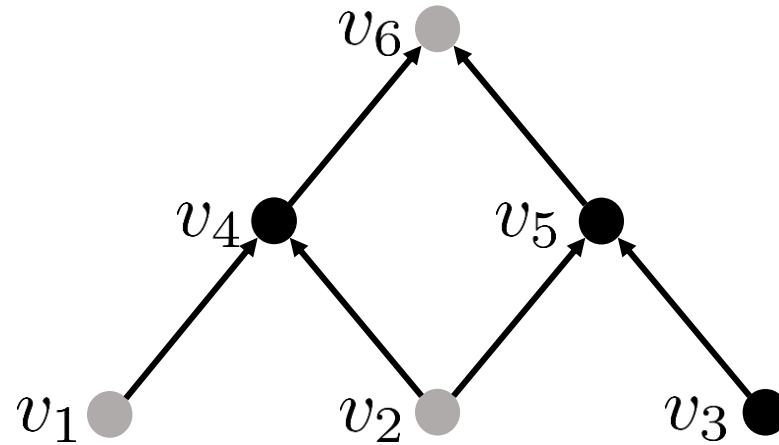


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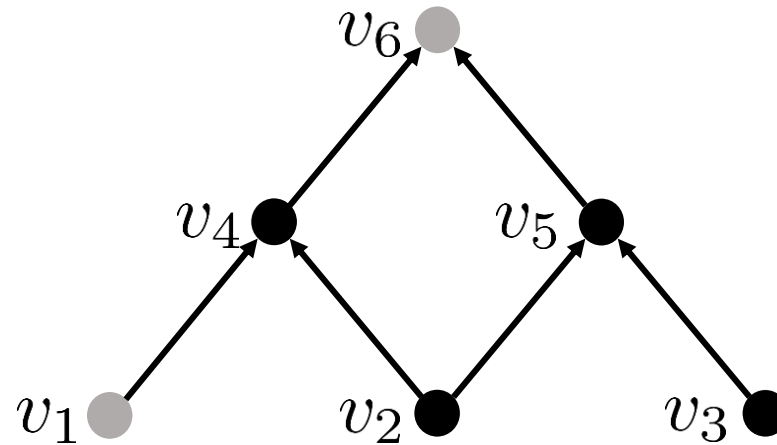


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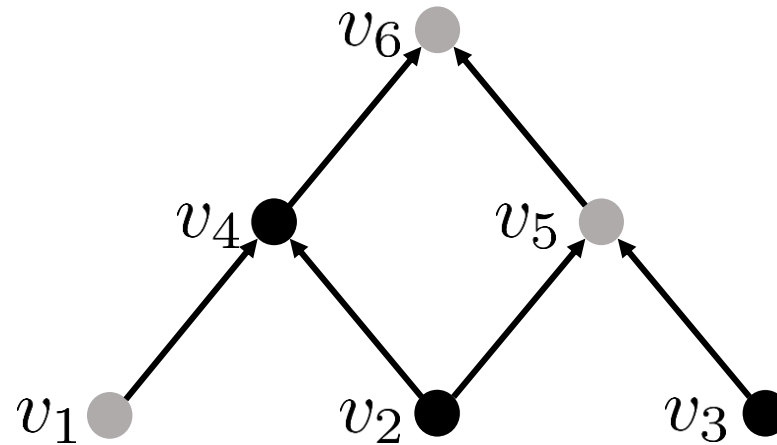
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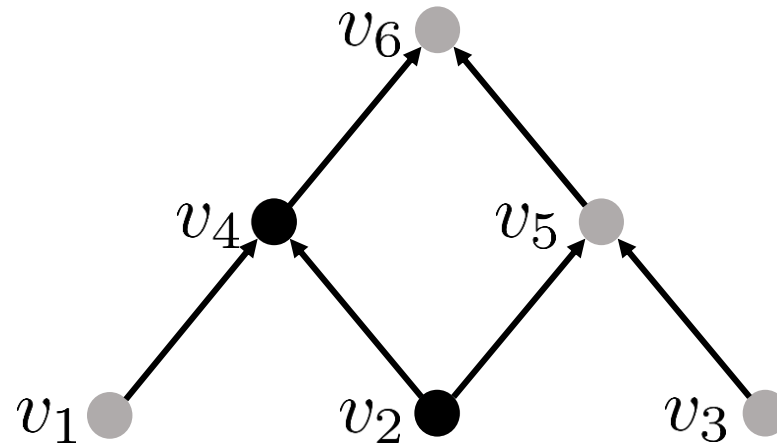


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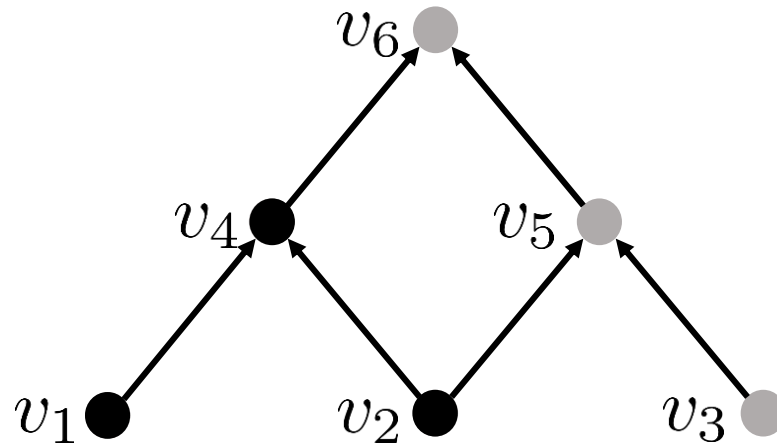


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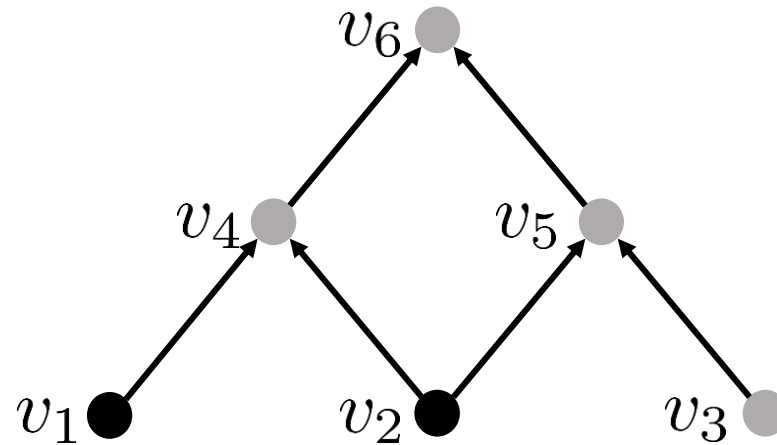


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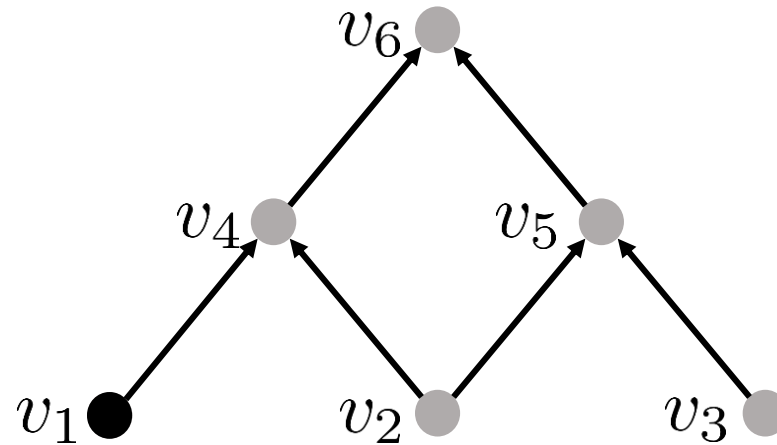


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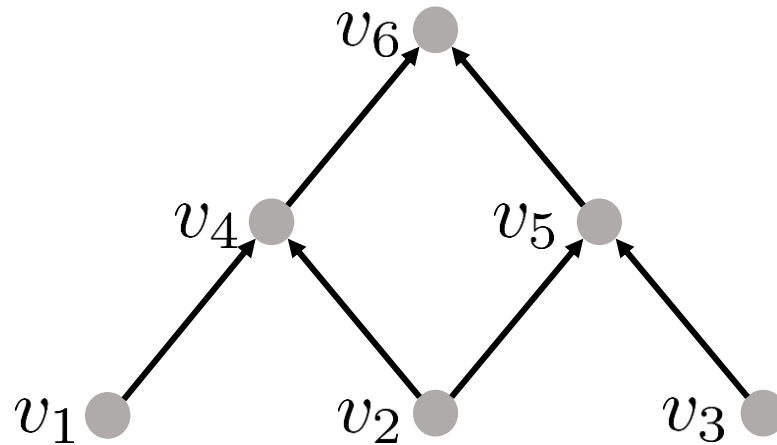


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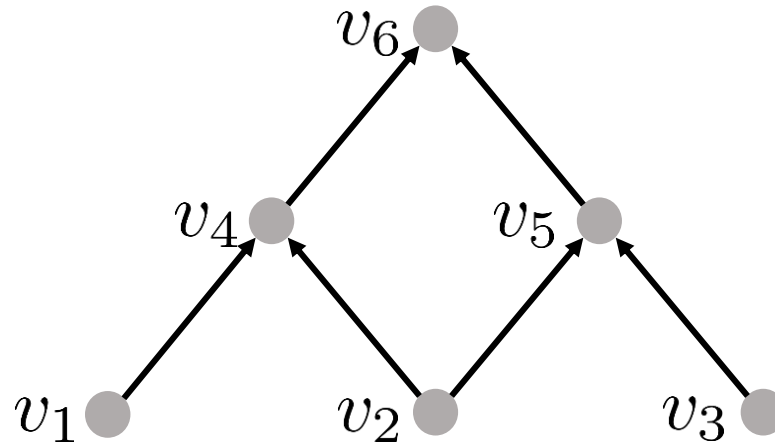
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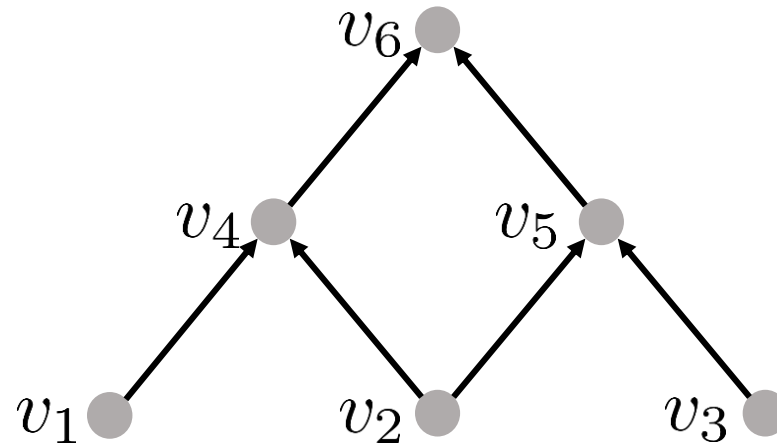
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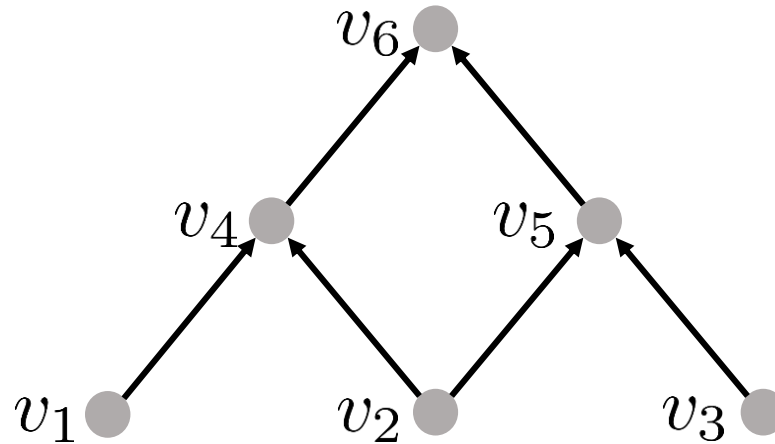
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- ▶ Space: maximum $\#$ pebbles in any configuration

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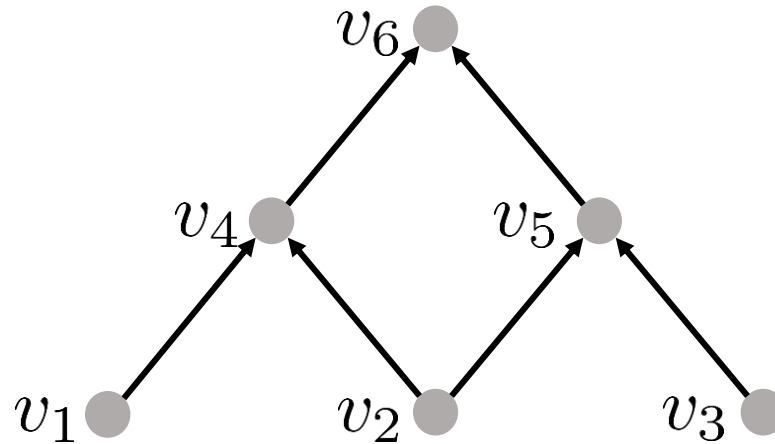
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- ▶ Space: maximum $\#$ pebbles in any configuration (4 in example)

Reversible pebble game



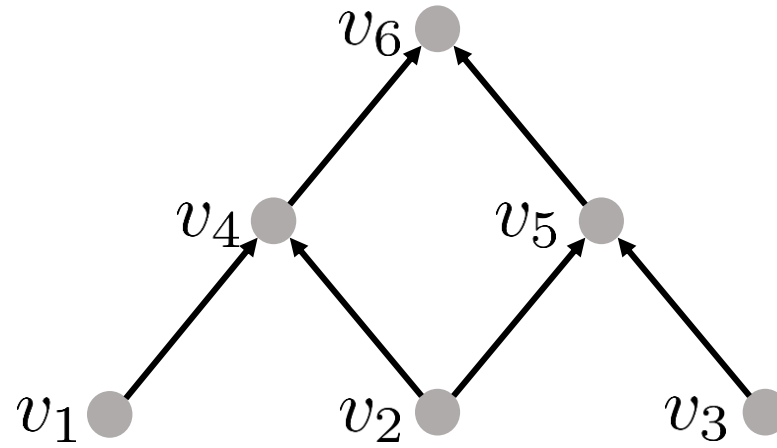
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- ▶ Time: $\#$ moves

Reversible pebble game



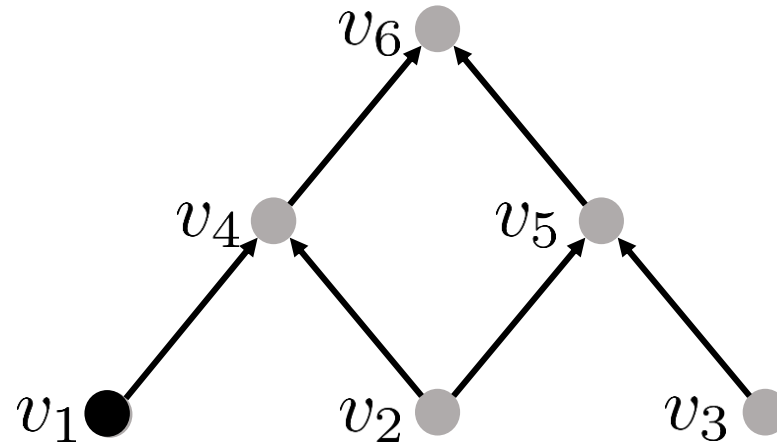
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- ▶ Time: $\#$ moves ($t = 16$ in example)

Reversible pebble game



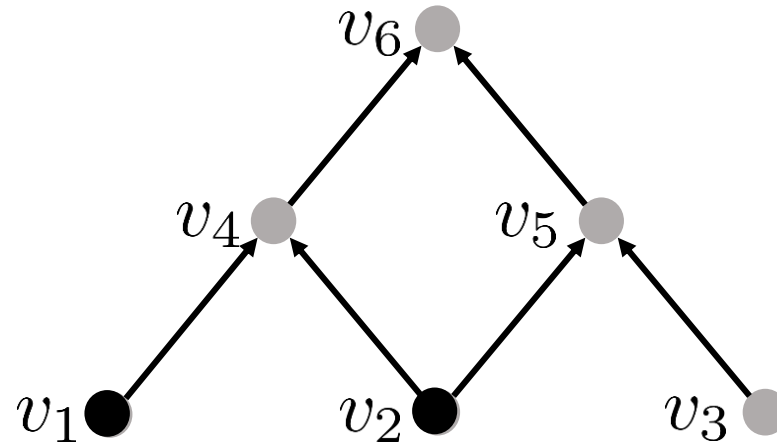
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- ▶ Faster pebbling?

Reversible pebble game



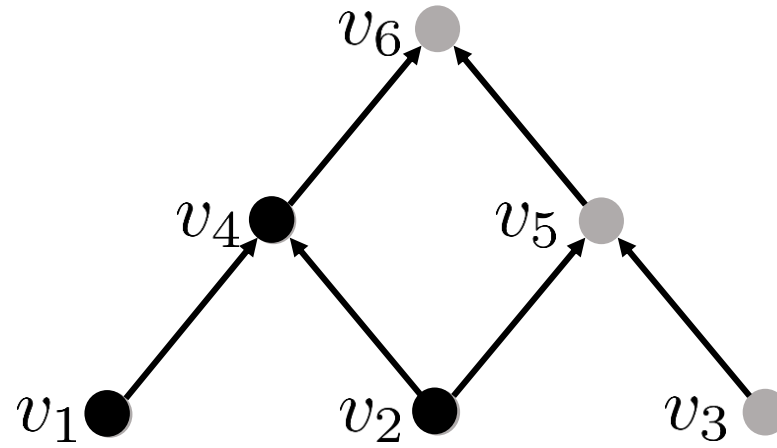
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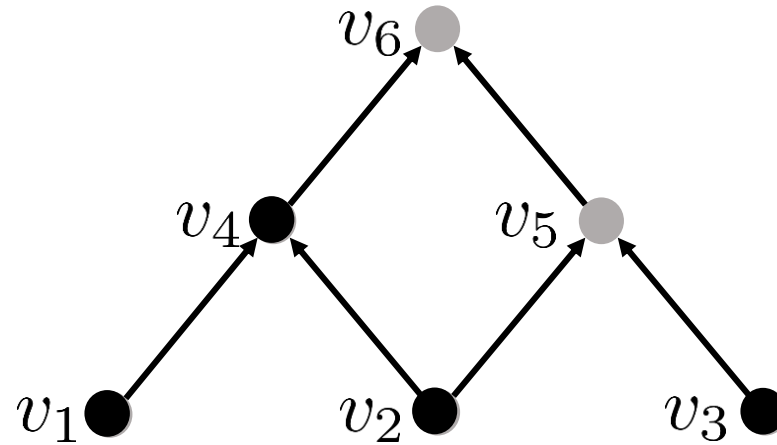
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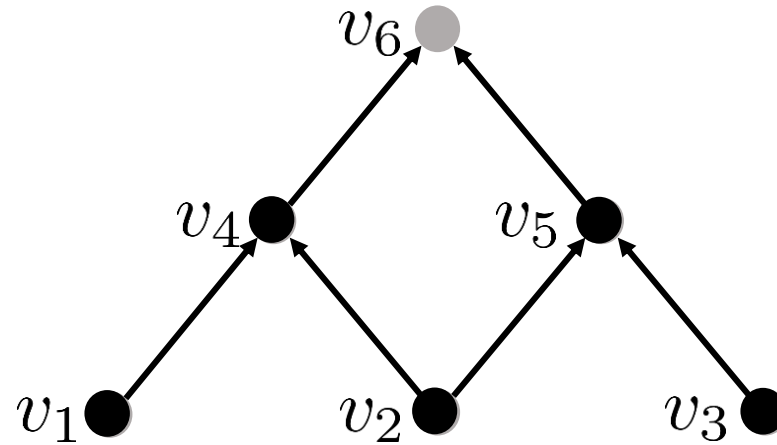
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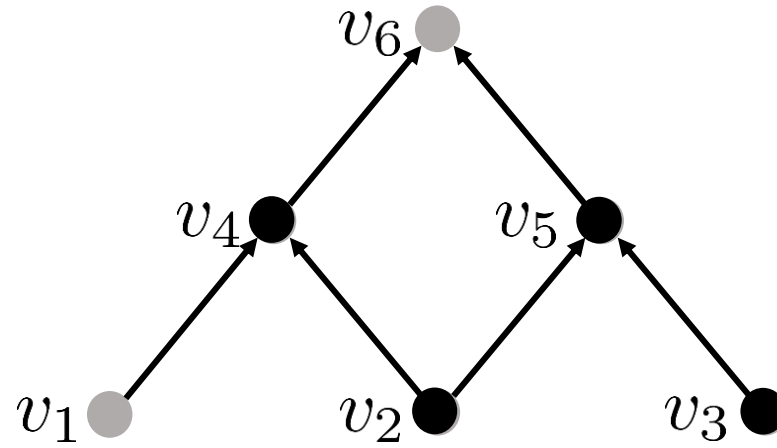
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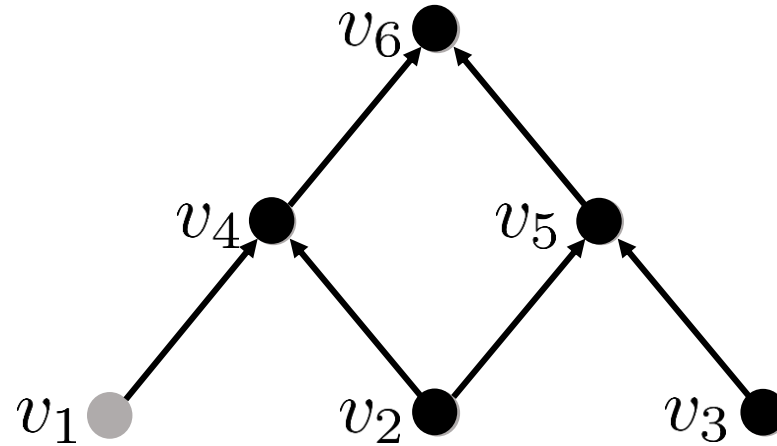
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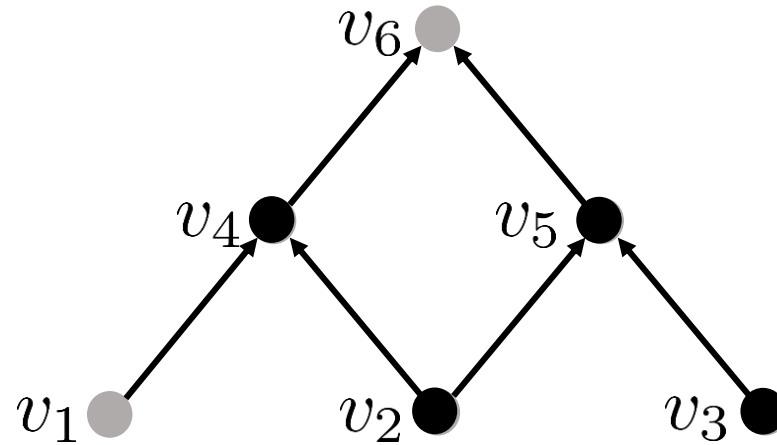
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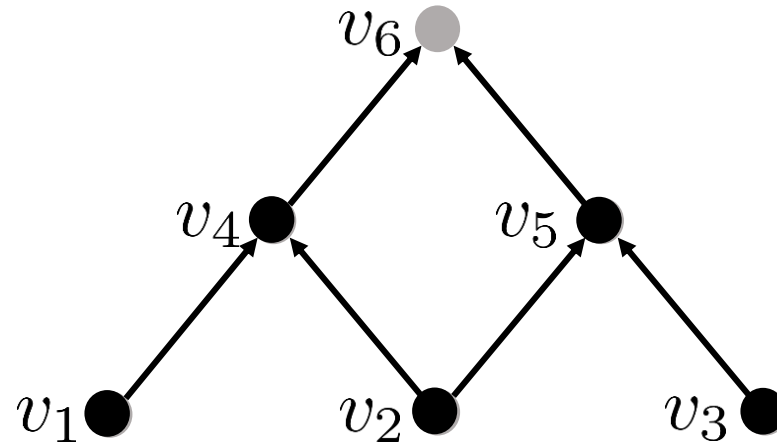
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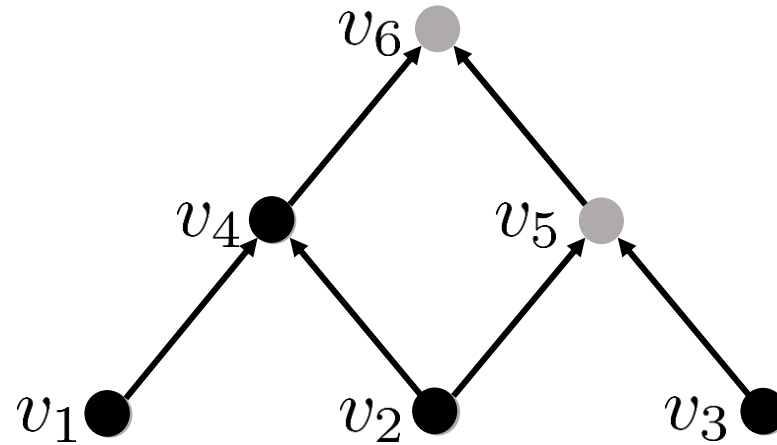
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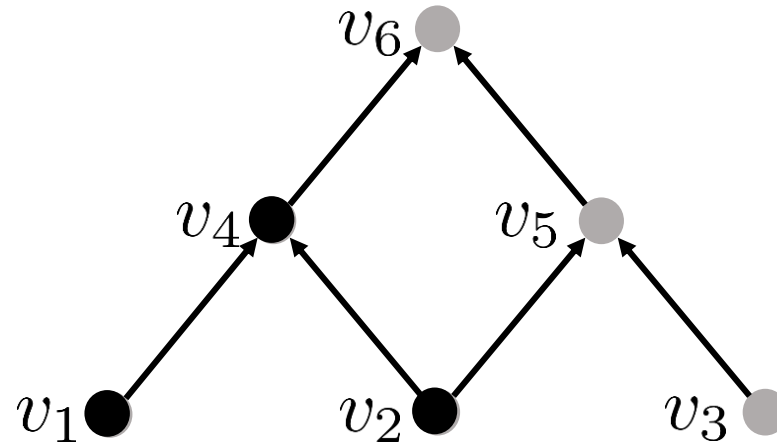
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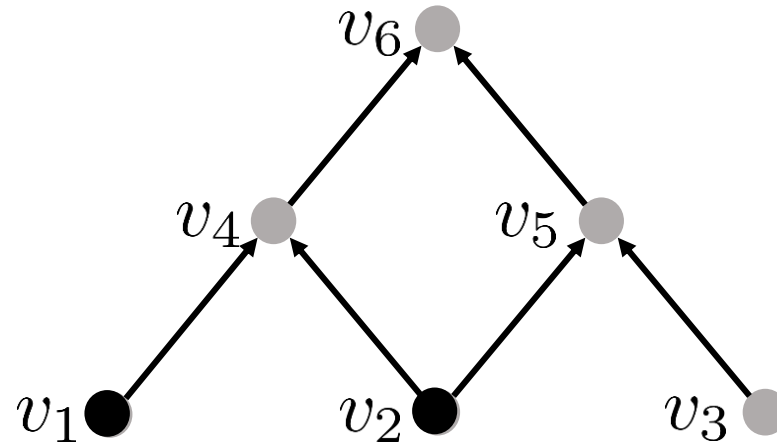
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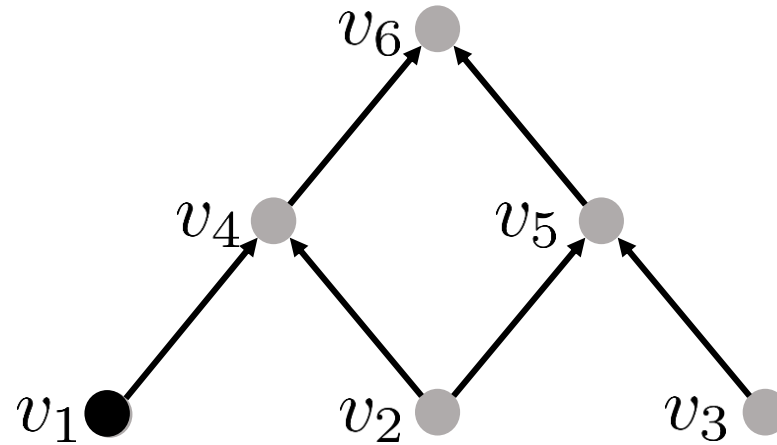
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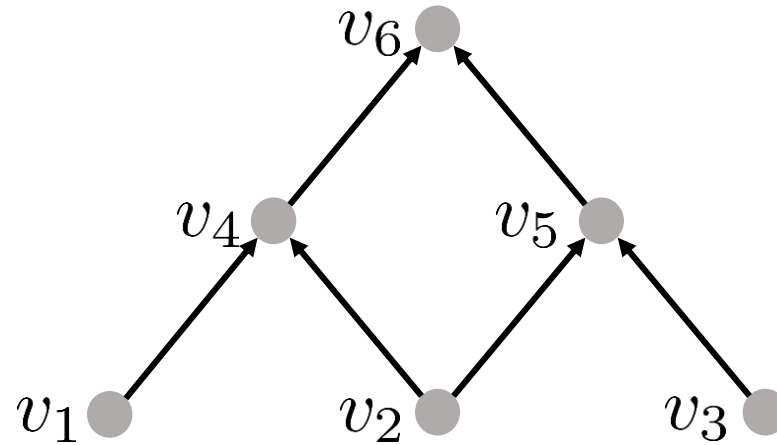
- ▶ Can do space 4, time 16
- ▶ Faster pebbling?

Reversible pebble game



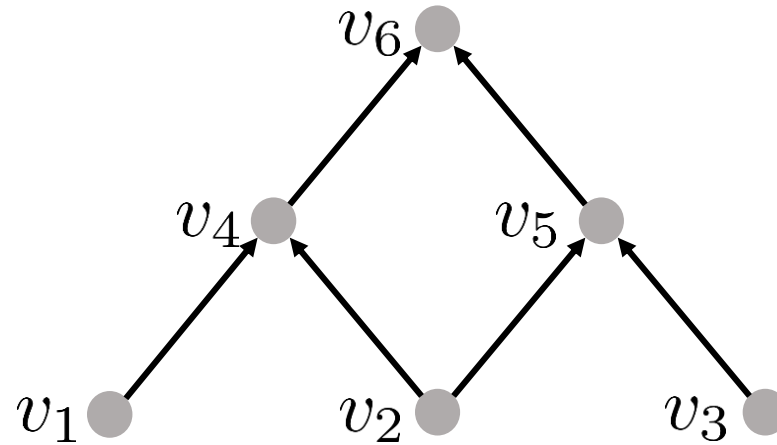
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Reversible pebble game



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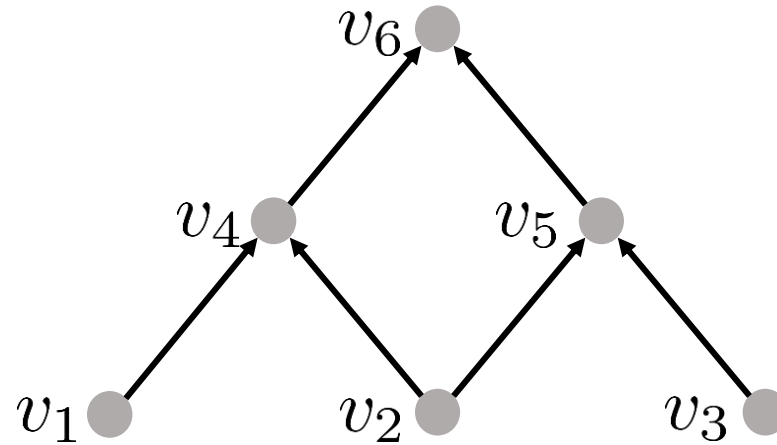
Reversible pebble game



- ▶ Can do space 4, time 16
- ▶ Faster pebbling?

space	time
4	16

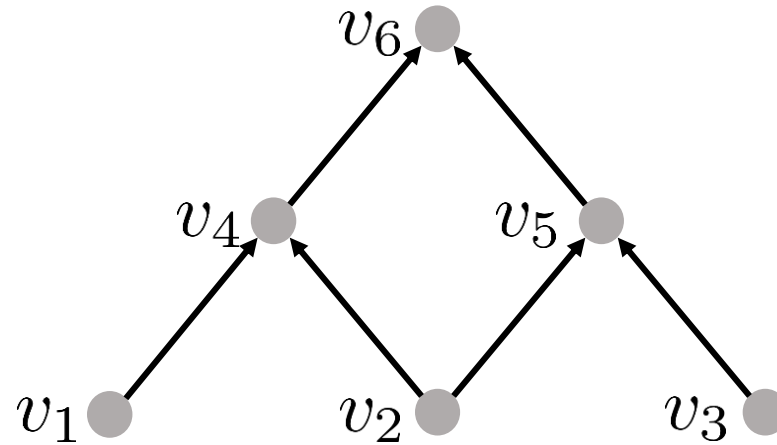
Reversible pebble game



- ▶ Can do space 4, time 16
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space	time
4	16
5	14

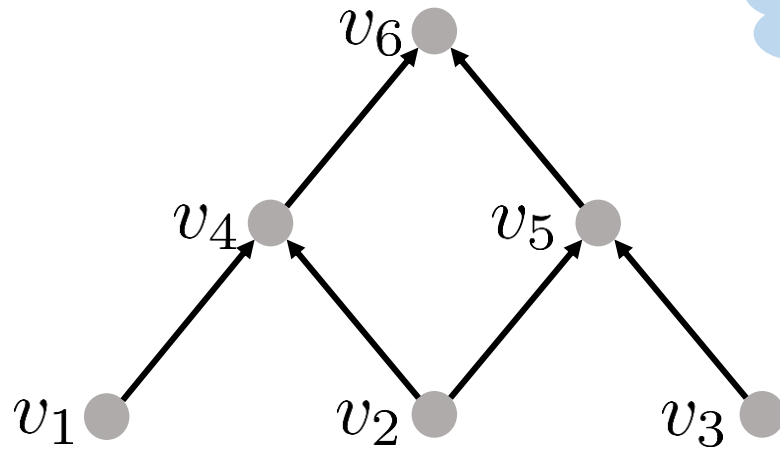
Reversible pebble game



- ▶ Can do space 4, time 16
- ▶ Faster pebbling?

space	time
4	16
5	14
6	12

Pebbling contradiction Peb_G



x true $\Leftrightarrow x = 1$
 x false $\Leftrightarrow x = 0$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

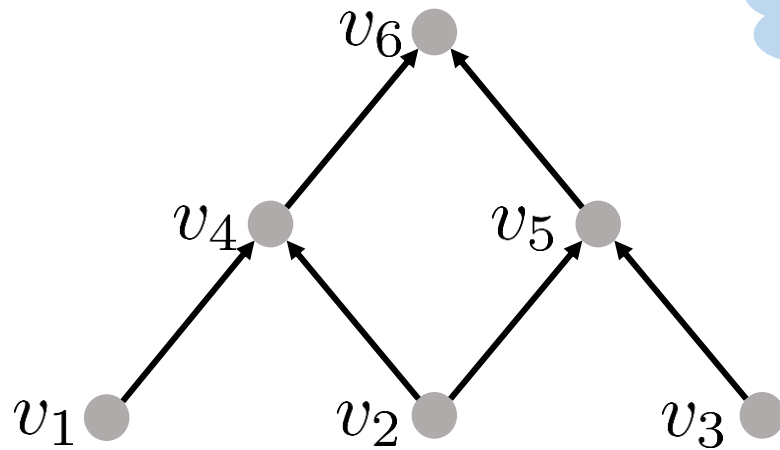
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

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$$x_{v_6} = 0$$

$$A_v := (1 - x_v) \prod_{u \in \text{pred}(v)} x_u, \text{ for all } v \in V(G)$$

$$A_{\text{sink}} := x_{\text{sink}}$$

$\text{pred}(v)$: set of all predecessors of v

Nullstellensatz refutation

- ▶ Pebbling contradiction Peb_G

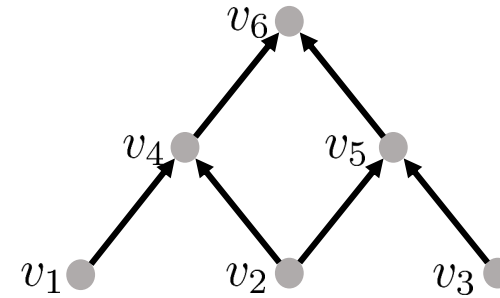
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- ▶ NS refutation

$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

Reversible pebbling

- ▶ DAG G



- ▶ Reversible pebbling

$$\emptyset = \mathbb{P}_0, \mathbb{P}_1, \dots, \mathbb{P}_t = \emptyset$$

Nullstellensatz refutation

► Pebbling contradiction Peb_G

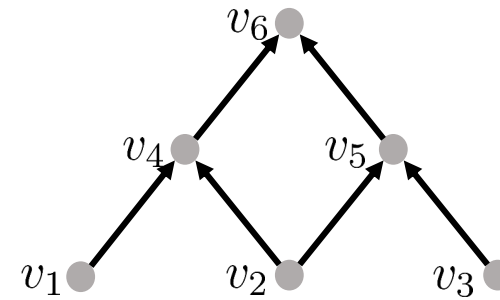
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Reversible pebbling

► DAG G



► Reversible pebbling

$$\emptyset = \mathbb{P}_0, \mathbb{P}_1, \dots, \mathbb{P}_t = \emptyset$$

Theorem

\exists NS refutation of Peb_G in size $t + 1$ and degree s

\Leftrightarrow

\exists reversible pebbling of G in time t and space s

Resolution

- ▶ Reason with clauses

$$\frac{y \vee \neg x \quad z \vee y \vee x}{z \vee y}$$

Resolution

- ▶ Reason with clauses
- ▶ Measure size, width

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Resolution

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$$\frac{y \vee \neg x \quad z \vee y \vee x}{z \vee y}$$

Polynomial calculus

- ▶ Reason with polynomials

$$\frac{2xy + 3xz = 0 \quad xyz - xz = 0}{2y + 3xz = 0}$$

Resolution

- ▶ Reason with clauses
- ▶ Measure size, width

$$\frac{y \vee \neg x \quad z \vee y \vee x}{z \vee y}$$

Polynomial calculus

- ▶ Reason with polynomials

$$\frac{2xy + 3xz = 0 \quad 3 \cdot (xyz - xz = 0)}{2xy + 3xyz = 0}$$

Resolution

- ▶ Reason with clauses
- ▶ Measure size, width

$$\frac{y \vee \neg x \quad z \vee y \vee x}{z \vee y}$$

Polynomial calculus

- ▶ Reason with polynomials
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Resolution

- ▶ Reason with clauses
- ▶ Measure size, width

$$\frac{y \vee \neg x \quad z \vee y \vee x}{z \vee y}$$

- ▶ Small degree/width \Rightarrow small size

Only $\binom{2n}{\leq w} \leq (2n)^w$ clauses of width $\leq w$ (essentially tight [ALN16])

Polynomial calculus

- ▶ Reason with polynomials
- ▶ Measure size, degree

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- ▶ Small size \Rightarrow (medium-)small degree/width [IPS99, BW01]

Polynomial calculus

- ▶ Reason with polynomials
- ▶ Measure size, degree

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Size-degree relation

Polynomial calculus, resolution

- ▶ Small degree/width \Rightarrow small size
- ▶ Small size \Rightarrow (medium-)small degree/width [IPS99, BW01]

Size-degree relation

Polynomial calculus, resolution

- ▶ Small degree/width \Rightarrow small size
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Nullstellensatz

- ▶ Small degree \Rightarrow small size
- ▶ Small size $\not\Rightarrow$ small degree

Size-degree relation

Polynomial calculus, resolution

- ▶ Small degree/width \Rightarrow small size
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Nullstellensatz

- ▶ Small degree \Rightarrow small size
- ▶ Small size $\not\Rightarrow$ small degree

Small size \Rightarrow small degree: reduction blows up size. Inherent?

- ▶ Resolution: yes, strong size-width trade-offs [Tha16]
- ▶ Polynomial calculus: open
- ▶ Nullstellensatz: strong size-degree trade-offs [this work]

Nullstellensatz size-degree trade-offs

Theorem

There is an explicit family of sets of polynomials s.t.

1. \exists NS refutation in nearly **linear size** and **degree d_1** ;

Nullstellensatz size-degree trade-offs

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1. \exists NS refutation in nearly **linear size** and **degree d_1** ;
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Nullstellensatz size-degree trade-offs

Theorem

There is an explicit family of sets of polynomials s.t.

1. \exists NS refutation in nearly **linear size** and **degree d_1** ;
2. \exists NS refutation in **degree $d_2 \ll d_1$** (and **size $\leq n^{d_2}$**);
3. any NS refutation in **degree slightly below d_1** has **size nearly n^{d_2}** .

Nullstellensatz size-degree trade-offs

Theorem

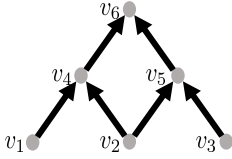
There is an explicit family of sets of polynomials s.t.

1. \exists NS refutation in nearly **linear size** and **degree d_1** ;
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Proof.

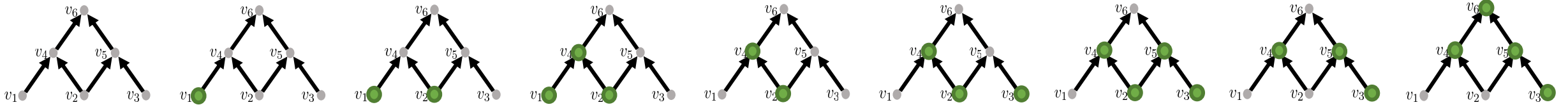
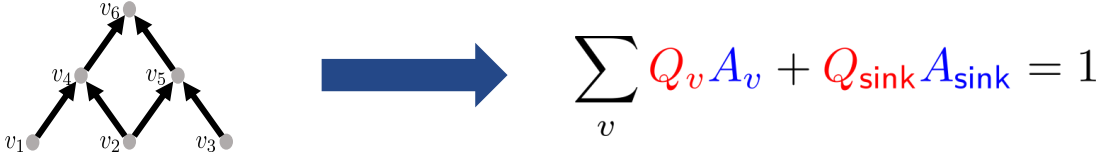
- ▶ \exists NS refutation in size $t + 1$, degree $s \Leftrightarrow \exists$ reversible pebbling in time t , space s
- ▶ Show strong reversible pebbling time-space trade-offs

Reversible pebbling to NS refutation

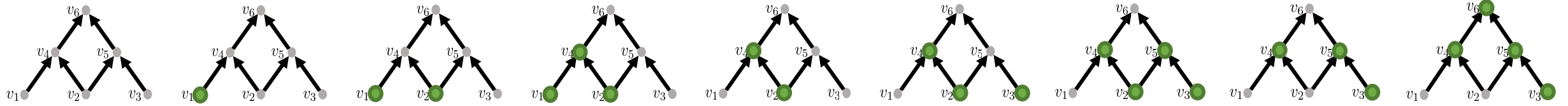


$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

Reversible pebbling to NS refutation



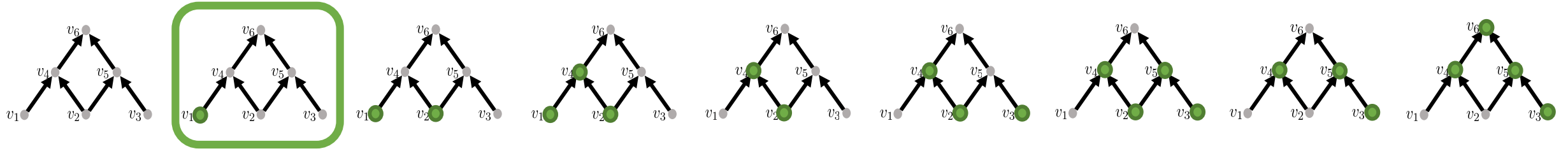
Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

$$\begin{aligned} 1 - x_{v_1} &= 0 \\ 1 - x_{v_2} &= 0 \\ 1 - x_{v_3} &= 0 \\ x_{v_1} x_{v_2} (1 - x_{v_4}) &= 0 \\ x_{v_2} x_{v_3} (1 - x_{v_5}) &= 0 \\ x_{v_4} x_{v_5} (1 - x_{v_6}) &= 0 \\ x_{v_6} &= 0 \end{aligned}$$

Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

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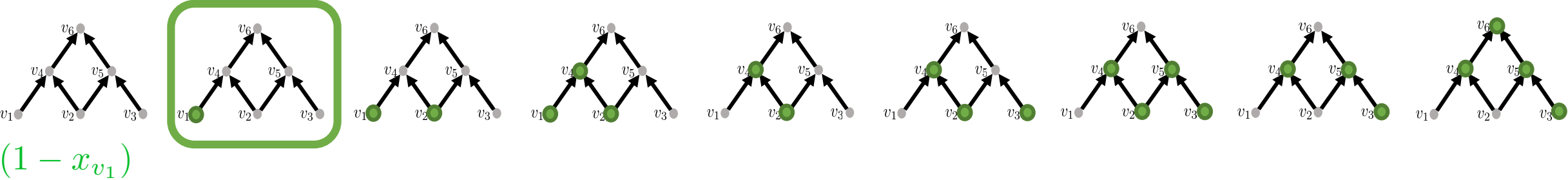
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

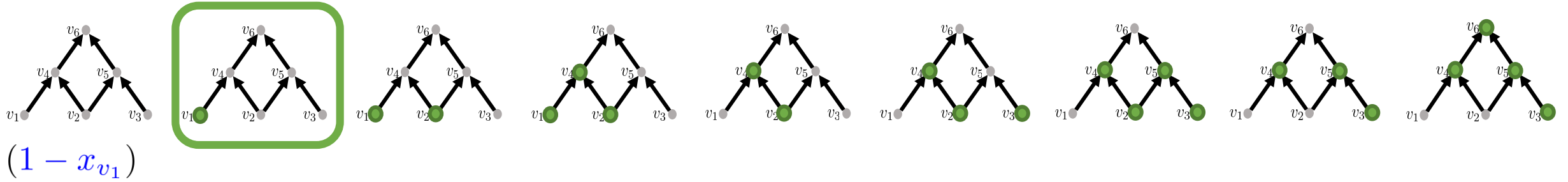
$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\begin{aligned}
 1 - x_{v_1} &= 0 \\
 1 - x_{v_2} &= 0 \\
 1 - x_{v_3} &= 0 \\
 x_{v_1} x_{v_2} (1 - x_{v_4}) &= 0 \\
 x_{v_2} x_{v_3} (1 - x_{v_5}) &= 0 \\
 x_{v_4} x_{v_5} (1 - x_{v_6}) &= 0 \\
 x_{v_6} &= 0
 \end{aligned}$$

Reversible pebbling to NS refutation



$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

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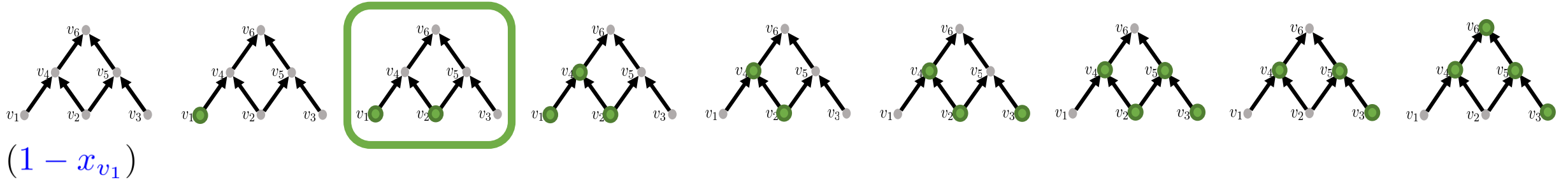
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$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

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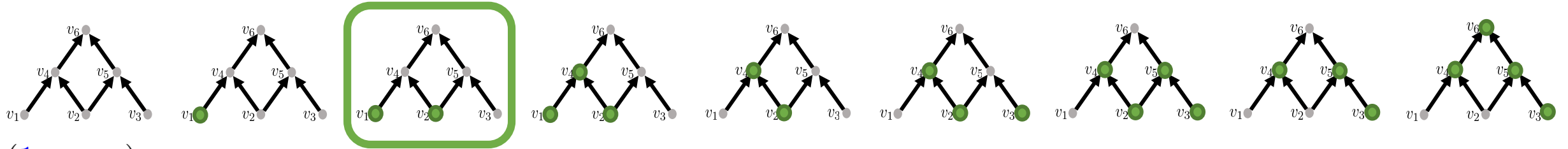
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

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$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

$$(1 - x_{v_1}) + x_{v_1}(1 - x_{v_2})$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

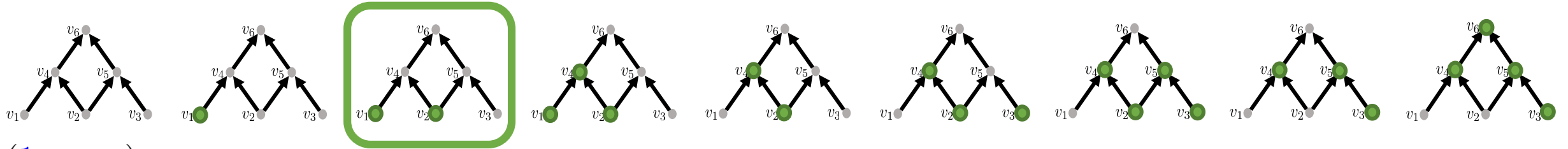
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$(1 - x_{v_1}) + x_{v_1}(1 - x_{v_2})$$

$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

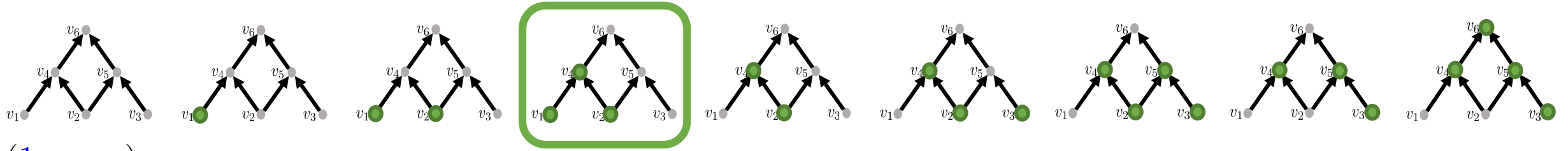
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\begin{aligned}
 &(1 - x_{v_1}) \\
 &+ x_{v_1}(1 - x_{v_2}) \\
 &+ x_{v_1}x_{v_2}(1 - x_{v_4})
 \end{aligned}$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

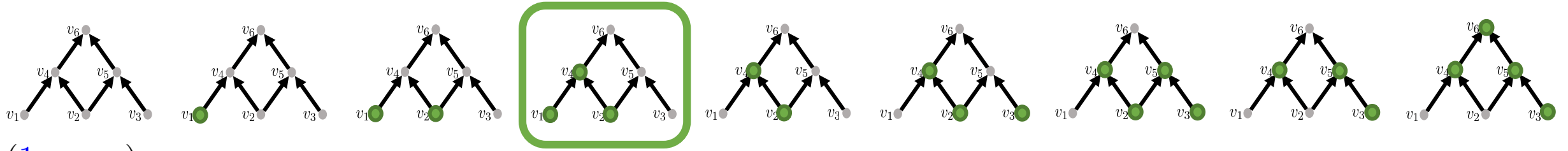
$$x_{v_1}x_{v_2}(1 - x_{v_4}) = 0$$

$$x_{v_2}x_{v_3}(1 - x_{v_5}) = 0$$

$$x_{v_4}x_{v_5}(1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

$$(1 - x_{v_1}) + x_{v_1}(1 - x_{v_2}) + x_{v_1}x_{v_2}(1 - x_{v_4})$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

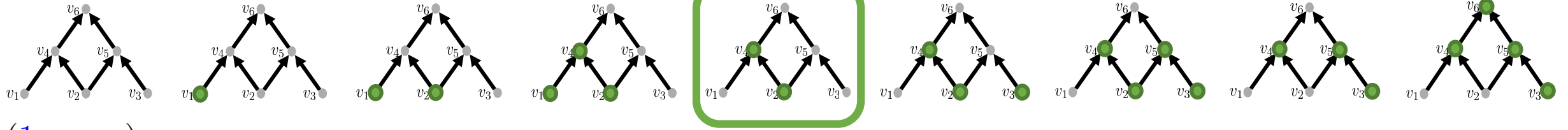
$$\rightarrow x_{v_1}x_{v_2}(1 - x_{v_4}) = 0$$

$$x_{v_2}x_{v_3}(1 - x_{v_5}) = 0$$

$$x_{v_4}x_{v_5}(1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\begin{aligned}
 &(1 - x_{v_1}) \\
 &+ x_{v_1}(1 - x_{v_2}) \\
 &\quad + x_{v_1}x_{v_2}(1 - x_{v_4}) \\
 &\quad\quad + x_{v_2}x_{v_4}(-1)(1 - x_{v_1})
 \end{aligned}$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

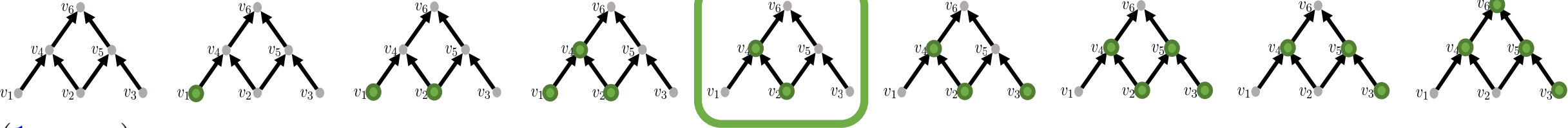
$$x_{v_1}x_{v_2}(1 - x_{v_4}) = 0$$

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Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

$$\begin{aligned} & (1 - x_{v_1}) \\ & + x_{v_1} (1 - x_{v_2}) \\ & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\ & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \end{aligned}$$



$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

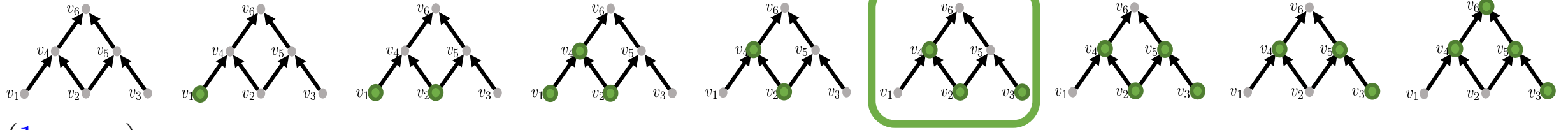
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

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Reversible pebbling to NS refutation



$$\begin{aligned}
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 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
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 \end{aligned}$$

$$1 - x_{v_1} = 0$$

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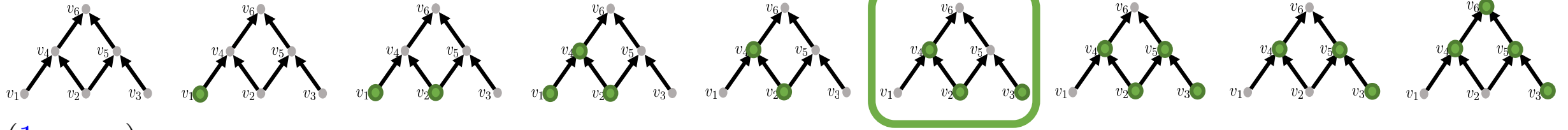
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

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$$1 - x_{v_1} = 0$$

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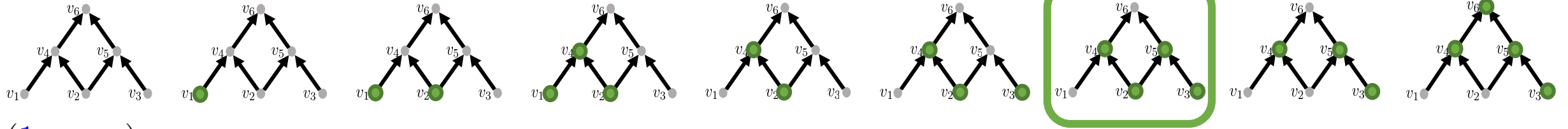
$$x_{v_1}x_{v_2}(1 - x_{v_4}) = 0$$

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Reversible pebbling to NS refutation



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 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5})
 \end{aligned}$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

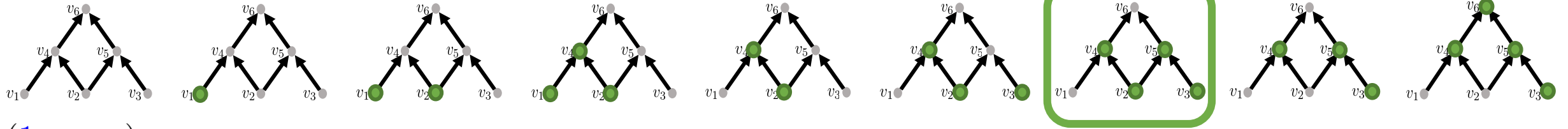
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

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Reversible pebbling to NS refutation



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$$\begin{aligned} & (1 - x_{v_1}) \\ & + x_{v_1} (1 - x_{v_2}) \\ & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\ & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\ & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\ & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \end{aligned}$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

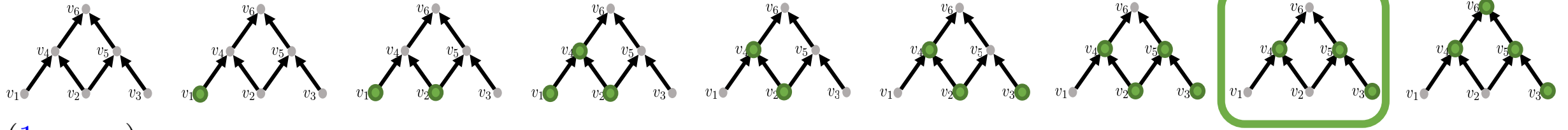
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2})
 \end{aligned}$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

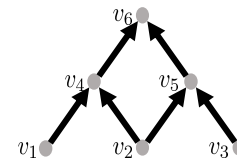
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

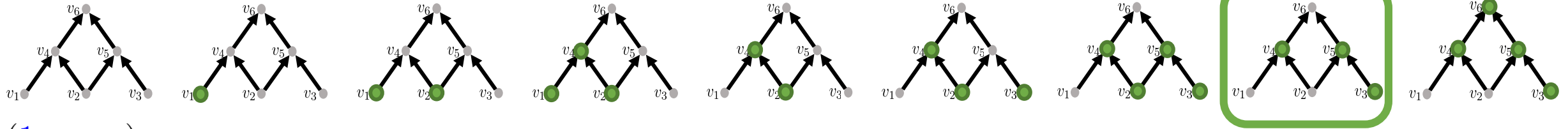
$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2})
 \end{aligned}$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

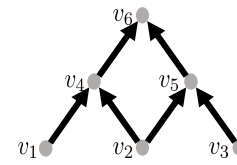
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

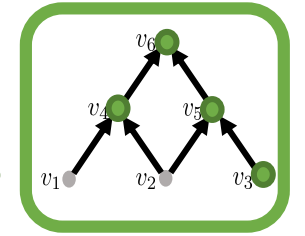
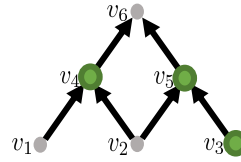
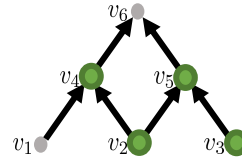
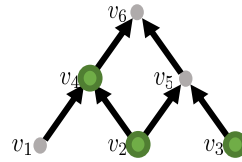
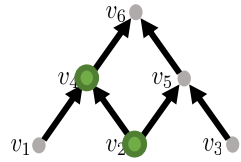
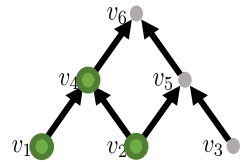
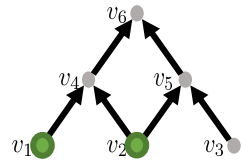
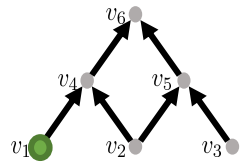
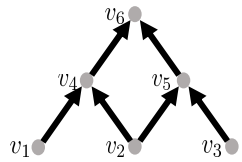
$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2}) \\
 & \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6})
 \end{aligned}$$

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

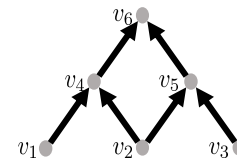
$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

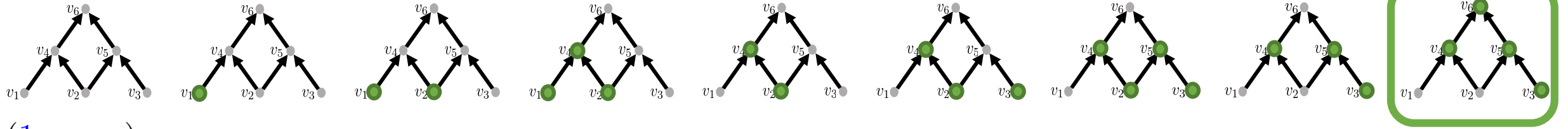
$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

Reversible pebbling to NS refutation

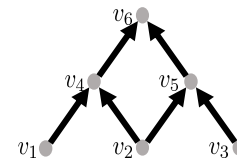


$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

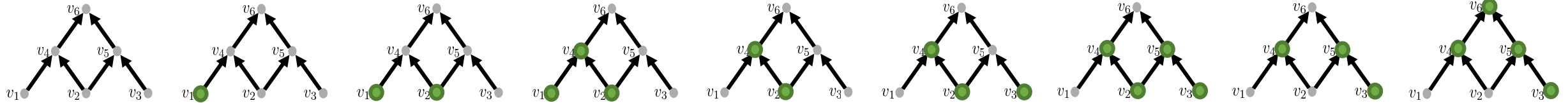


$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2}) \\
 & \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6}) \\
 & 1 - x_{v_1} = 0 \\
 & 1 - x_{v_2} = 0 \\
 & 1 - x_{v_3} = 0 \\
 & x_{v_1} x_{v_2} (1 - x_{v_4}) = 0 \\
 & x_{v_2} x_{v_3} (1 - x_{v_5}) = 0 \\
 \rightarrow & x_{v_4} x_{v_5} (1 - x_{v_6}) = 0 \\
 & x_{v_6} = 0
 \end{aligned}$$

Reversible pebbling to NS refutation

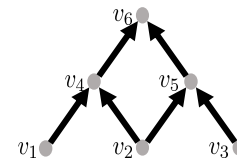


$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

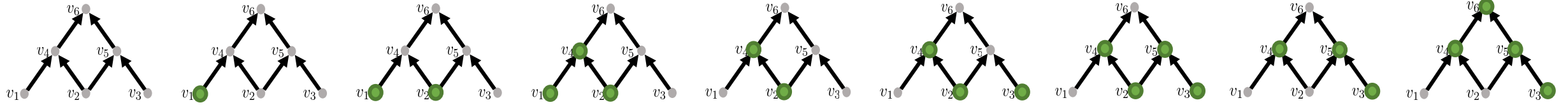


$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2}) \\
 & \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6}) \\
 & 1 - x_{v_1} = 0 \\
 & 1 - x_{v_2} = 0 \\
 & 1 - x_{v_3} = 0 \\
 & x_{v_1} x_{v_2} (1 - x_{v_4}) = 0 \\
 & x_{v_2} x_{v_3} (1 - x_{v_5}) = 0 \\
 & x_{v_4} x_{v_5} (1 - x_{v_6}) = 0 \\
 & x_{v_6} = 0
 \end{aligned}$$

Reversible pebbling to NS refutation

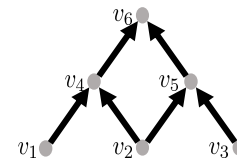


$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

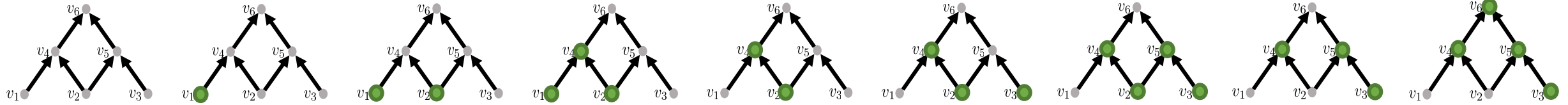


$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + (x_{v_1} - x_{v_1} x_{v_2}) \\
 & \quad + (x_{v_1} x_{v_2} - x_{v_1} x_{v_2} x_{v_4}) \\
 & \quad \quad + (x_{v_1} x_{v_2} x_{v_4} - x_{v_2} x_{v_4}) \\
 & \quad \quad \quad + (x_{v_2} x_{v_4} - x_{v_2} x_{v_3} x_{v_4}) \\
 & \quad \quad \quad \quad + (x_{v_2} x_{v_3} x_{v_4} - x_{v_2} x_{v_3} x_{v_4} x_{v_5}) \\
 & \quad \quad \quad \quad \quad + (x_{v_2} x_{v_3} x_{v_4} x_{v_5} - x_{v_3} x_{v_4} x_{v_5}) \\
 & \quad \quad \quad \quad \quad \quad + (x_{v_3} x_{v_4} x_{v_5} - x_{v_3} x_{v_4} x_{v_5} x_{v_6}) \\
 & 1 - x_{v_1} = 0 \\
 & 1 - x_{v_2} = 0 \\
 & 1 - x_{v_3} = 0 \\
 & x_{v_1} x_{v_2} (1 - x_{v_4}) = 0 \\
 & x_{v_2} x_{v_3} (1 - x_{v_5}) = 0 \\
 & x_{v_4} x_{v_5} (1 - x_{v_6}) = 0 \\
 & x_{v_6} = 0
 \end{aligned}$$

Reversible pebbling to NS refutation

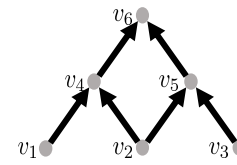


$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

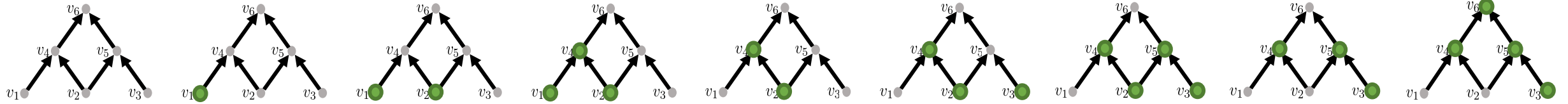


$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + (x_{v_1} - x_{v_1} x_{v_2}) \\
 & \quad + (x_{v_1} x_{v_2} - x_{v_1} x_{v_2} x_{v_4}) \\
 & \quad \quad + (x_{v_1} x_{v_2} x_{v_4} - x_{v_2} x_{v_4}) \\
 & \quad \quad \quad + (x_{v_2} x_{v_4} - x_{v_2} x_{v_3} x_{v_4}) \\
 & \quad \quad \quad \quad + (x_{v_2} x_{v_3} x_{v_4} - x_{v_2} x_{v_3} x_{v_4} x_{v_5}) \\
 & \quad \quad \quad \quad \quad + (x_{v_2} x_{v_3} x_{v_4} x_{v_5} - x_{v_3} x_{v_4} x_{v_5}) \\
 & \quad \quad \quad \quad \quad \quad + (x_{v_3} x_{v_4} x_{v_5} - x_{v_3} x_{v_4} x_{v_5} x_{v_6}) \\
 & \quad \quad \quad \quad \quad \quad \quad = 1 - x_{v_3} x_{v_4} x_{v_5} x_{v_6} \\
 \\
 & 1 - x_{v_1} = 0 \\
 & 1 - x_{v_2} = 0 \\
 & 1 - x_{v_3} = 0 \\
 & x_{v_1} x_{v_2} (1 - x_{v_4}) = 0 \\
 & x_{v_2} x_{v_3} (1 - x_{v_5}) = 0 \\
 & x_{v_4} x_{v_5} (1 - x_{v_6}) = 0 \\
 & x_{v_6} = 0
 \end{aligned}$$

Reversible pebbling to NS refutation



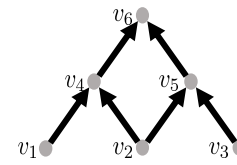
$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



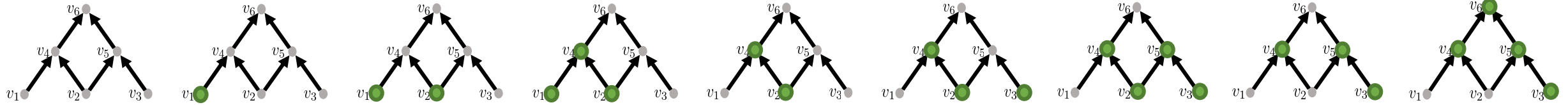
$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2}) \\
 & \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6}) \\
 \\
 & 1 - x_{v_1} = 0 \\
 & 1 - x_{v_2} = 0 \\
 & 1 - x_{v_3} = 0 \\
 & x_{v_1} x_{v_2} (1 - x_{v_4}) = 0 \\
 & x_{v_2} x_{v_3} (1 - x_{v_5}) = 0 \\
 & x_{v_4} x_{v_5} (1 - x_{v_6}) = 0 \\
 & x_{v_6} = 0
 \end{aligned}$$

$$= 1 - x_{v_3} x_{v_4} x_{v_5} x_{v_6}$$

Reversible pebbling to NS refutation



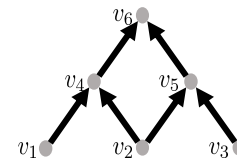
$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



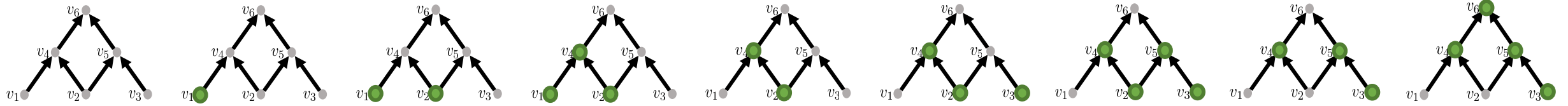
$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2}) \\
 & \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6}) \\
 & \quad \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} x_{v_6} \\
 & \quad \quad \quad \quad \quad \quad \quad \quad = 1
 \end{aligned}$$

$$\begin{aligned}
 1 - x_{v_1} &= 0 \\
 1 - x_{v_2} &= 0 \\
 1 - x_{v_3} &= 0 \\
 x_{v_1} x_{v_2} (1 - x_{v_4}) &= 0 \\
 x_{v_2} x_{v_3} (1 - x_{v_5}) &= 0 \\
 x_{v_4} x_{v_5} (1 - x_{v_6}) &= 0 \\
 x_{v_6} &= 0
 \end{aligned}$$

Reversible pebbling to NS refutation



$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$

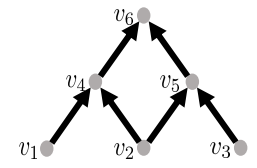


$$\begin{aligned}
 & (1 - x_{v_1}) \\
 & + x_{v_1} (1 - x_{v_2}) \\
 & \quad + x_{v_1} x_{v_2} (1 - x_{v_4}) \\
 & \quad \quad + x_{v_2} x_{v_4} (-1) (1 - x_{v_1}) \\
 & \quad \quad \quad + x_{v_2} x_{v_4} (1 - x_{v_3}) \\
 & \quad \quad \quad \quad + x_{v_2} x_{v_3} x_{v_4} (1 - x_{v_5}) \\
 & \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (-1) (1 - x_{v_2}) \\
 & \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6}) \\
 & \quad \quad \quad \quad \quad \quad \quad + x_{v_3} x_{v_4} x_{v_5} x_{v_6} \\
 & \quad \quad \quad \quad \quad \quad \quad \quad = 1
 \end{aligned}$$

NS size = $2 \cdot \frac{\text{pebbling time}}{2} + 1 = 2 \cdot 8 + 1$
 NS degree = pebbling space = 4

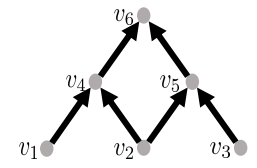
NS refutation to reversible pebbling

$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



NS refutation to reversible pebbling

$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

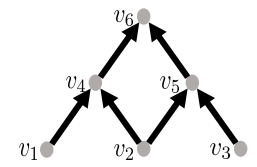
$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

$$\begin{aligned} (1 - x_{v_2} x_{v_4})(1 - x_{v_1}) + (x_{v_1} - x_{v_3} x_{v_4} x_{v_5})(1 - x_{v_2}) + x_{v_2} x_{v_4} (1 - x_{v_3}) + x_{v_1} x_{v_2} (1 - x_{v_4}) \\ + x_{v_4} x_{v_2} x_{v_3} (1 - x_{v_5}) + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6}) + x_{v_3} x_{v_4} x_{v_5} x_{v_6} = 1 \end{aligned}$$

NS refutation to reversible pebbling

$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



For simplicity, assume \mathbb{F}_2

$$1 - x_{v_1} = 0$$

$$1 - x_{v_2} = 0$$

$$1 - x_{v_3} = 0$$

$$x_{v_1} x_{v_2} (1 - x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 - x_{v_5}) = 0$$

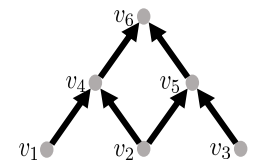
$$x_{v_4} x_{v_5} (1 - x_{v_6}) = 0$$

$$x_{v_6} = 0$$

$$\begin{aligned} (1 - x_{v_2} x_{v_4})(1 - x_{v_1}) + (x_{v_1} - x_{v_3} x_{v_4} x_{v_5})(1 - x_{v_2}) + x_{v_2} x_{v_4} (1 - x_{v_3}) + x_{v_1} x_{v_2} (1 - x_{v_4}) \\ + x_{v_4} x_{v_2} x_{v_3} (1 - x_{v_5}) + x_{v_3} x_{v_4} x_{v_5} (1 - x_{v_6}) + x_{v_3} x_{v_4} x_{v_5} x_{v_6} = 1 \end{aligned}$$

NS refutation to reversible pebbling

$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



For simplicity, assume \mathbb{F}_2

$$1 + x_{v_1} = 0$$

$$1 + x_{v_2} = 0$$

$$1 + x_{v_3} = 0$$

$$x_{v_1} x_{v_2} (1 + x_{v_4}) = 0$$

$$x_{v_2} x_{v_3} (1 + x_{v_5}) = 0$$

$$x_{v_4} x_{v_5} (1 + x_{v_6}) = 0$$

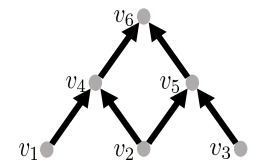
$$x_{v_6} = 0$$

$$\begin{aligned} (1 + x_{v_2} x_{v_4})(1 + x_{v_1}) + (x_{v_1} + x_{v_3} x_{v_4} x_{v_5})(1 + x_{v_2}) + x_{v_2} x_{v_4} (1 + x_{v_3}) + x_{v_1} x_{v_2} (1 + x_{v_4}) \\ + x_{v_4} x_{v_2} x_{v_3} (1 + x_{v_5}) + x_{v_3} x_{v_4} x_{v_5} (1 + x_{v_6}) + x_{v_3} x_{v_4} x_{v_5} x_{v_6} = 1 \end{aligned}$$

NS refutation to reversible pebbling

\forall monomials x_W in proof, add node W

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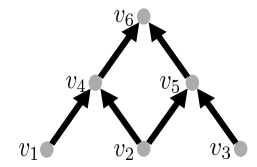
$$\begin{aligned} & (1 + x_{v_2} x_{v_4})(1 + x_{v_1}) + (x_{v_1} + x_{v_3} x_{v_4} x_{v_5})(1 + x_{v_2}) + x_{v_2} x_{v_4} (1 + x_{v_3}) + x_{v_1} x_{v_2} (1 + x_{v_4}) \\ & + x_{v_4} x_{v_2} x_{v_3} (1 + x_{v_5}) + x_{v_3} x_{v_4} x_{v_5} (1 + x_{v_6}) + x_{v_3} x_{v_4} x_{v_5} x_{v_6} = 1 \end{aligned}$$

NS refutation to reversible pebbling

\forall monomials x_W in proof, add node W

$$x_W = \prod_{v \in W} x_v$$

$$\sum_v Q_v A_v + Q_{\text{sink}} A_{\text{sink}} = 1$$



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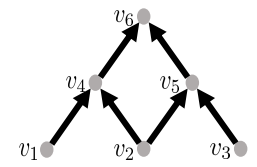
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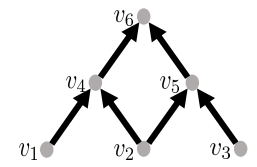
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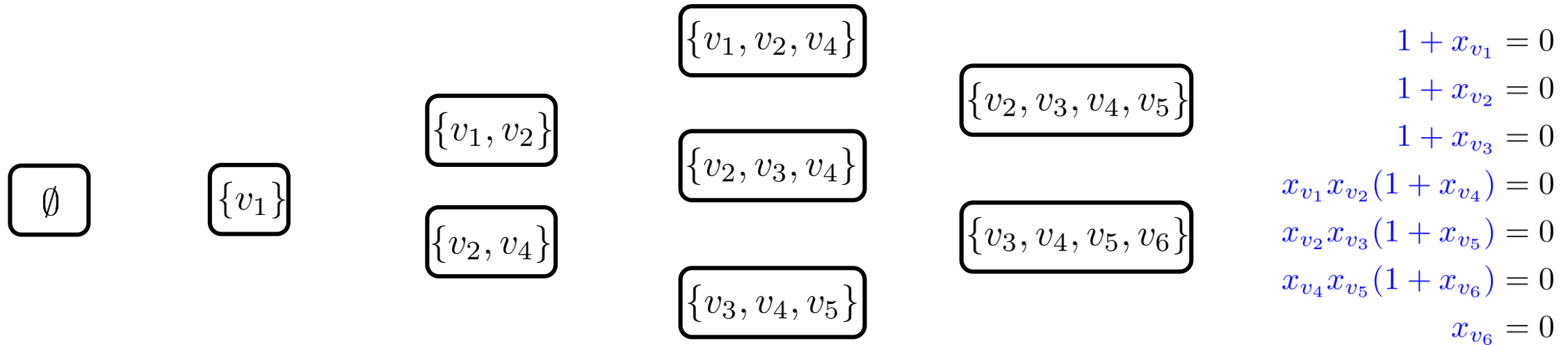
NS refutation to reversible pebbling

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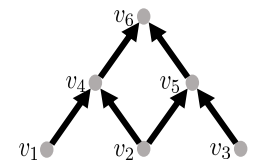
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NS refutation to reversible pebbling

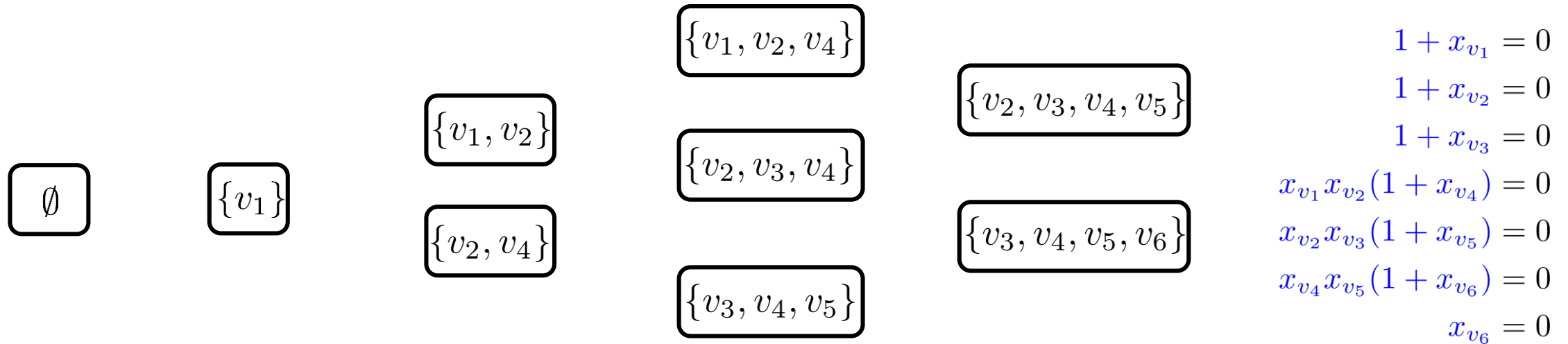
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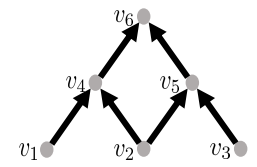
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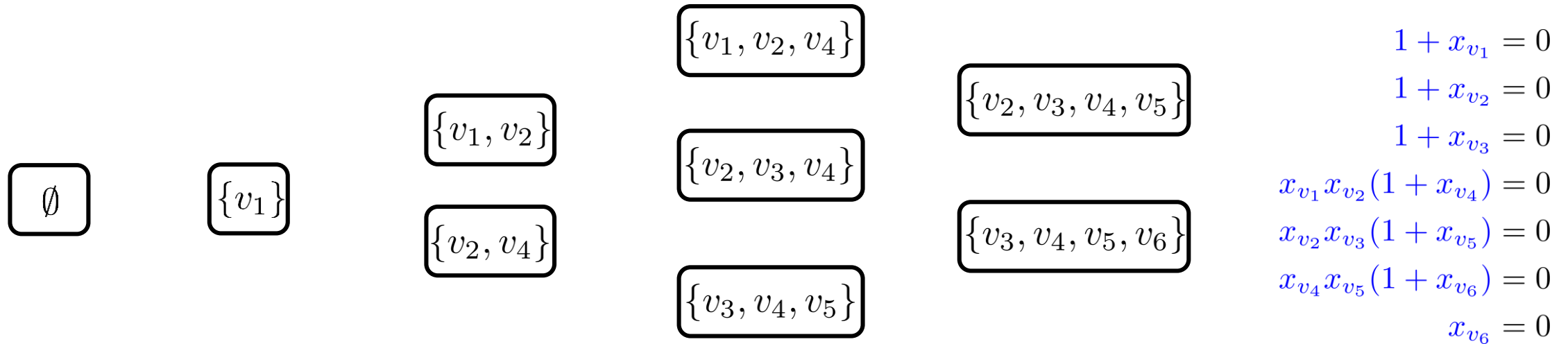
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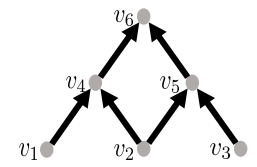
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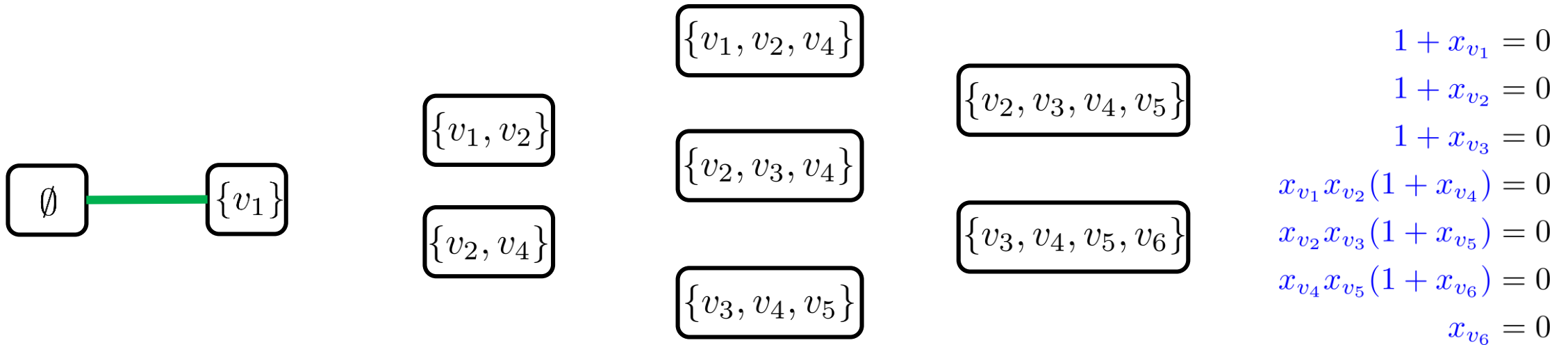
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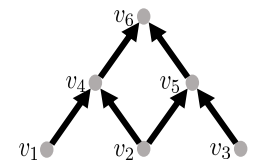
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$$\begin{aligned} & \underline{(1 + x_{v_2} x_{v_4})} \underline{(1 + x_{v_1})} + (x_{v_1} + x_{v_3} x_{v_4} x_{v_5})(1 + x_{v_2}) + x_{v_2} x_{v_4} (1 + x_{v_3}) + x_{v_1} x_{v_2} (1 + x_{v_4}) \\ & + x_{v_4} x_{v_2} x_{v_3} (1 + x_{v_5}) + x_{v_3} x_{v_4} x_{v_5} (1 + x_{v_6}) + x_{v_3} x_{v_4} x_{v_5} x_{v_6} = 1 \end{aligned}$$

NS refutation to reversible pebbling

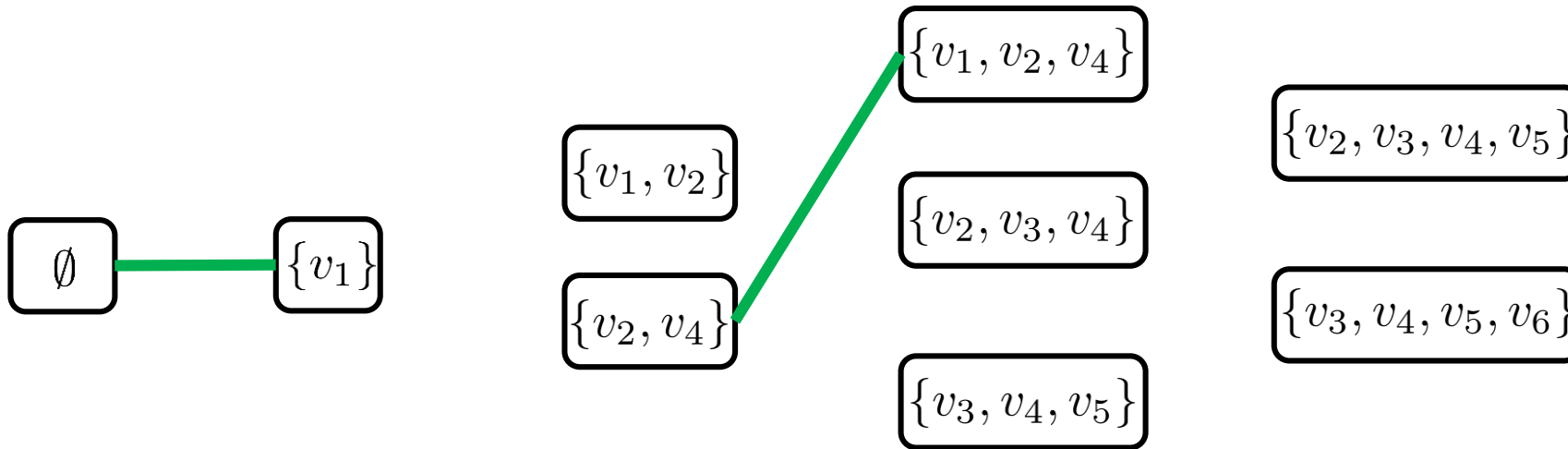
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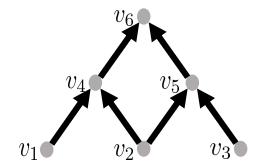


$$\begin{aligned} 1 + x_{v_1} &= 0 \\ 1 + x_{v_2} &= 0 \\ 1 + x_{v_3} &= 0 \\ x_{v_1} x_{v_2} (1 + x_{v_4}) &= 0 \\ x_{v_2} x_{v_3} (1 + x_{v_5}) &= 0 \\ x_{v_4} x_{v_5} (1 + x_{v_6}) &= 0 \\ x_{v_6} &= 0 \end{aligned}$$

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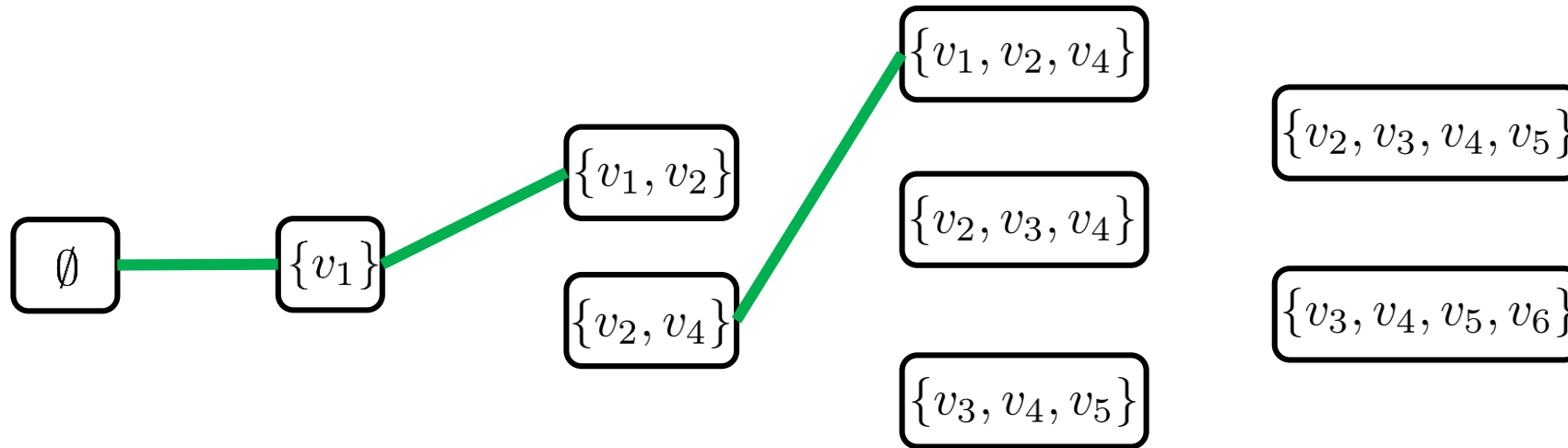
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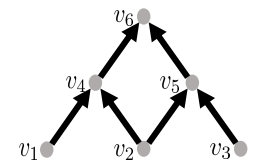


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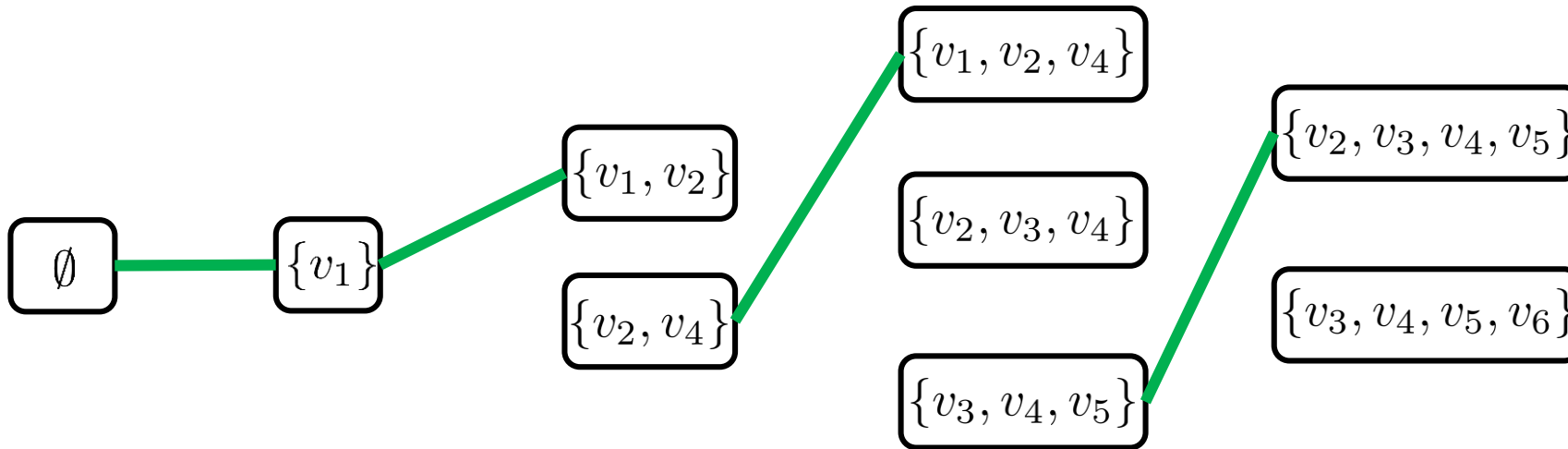
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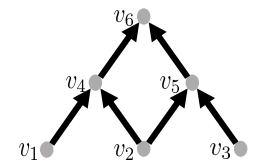


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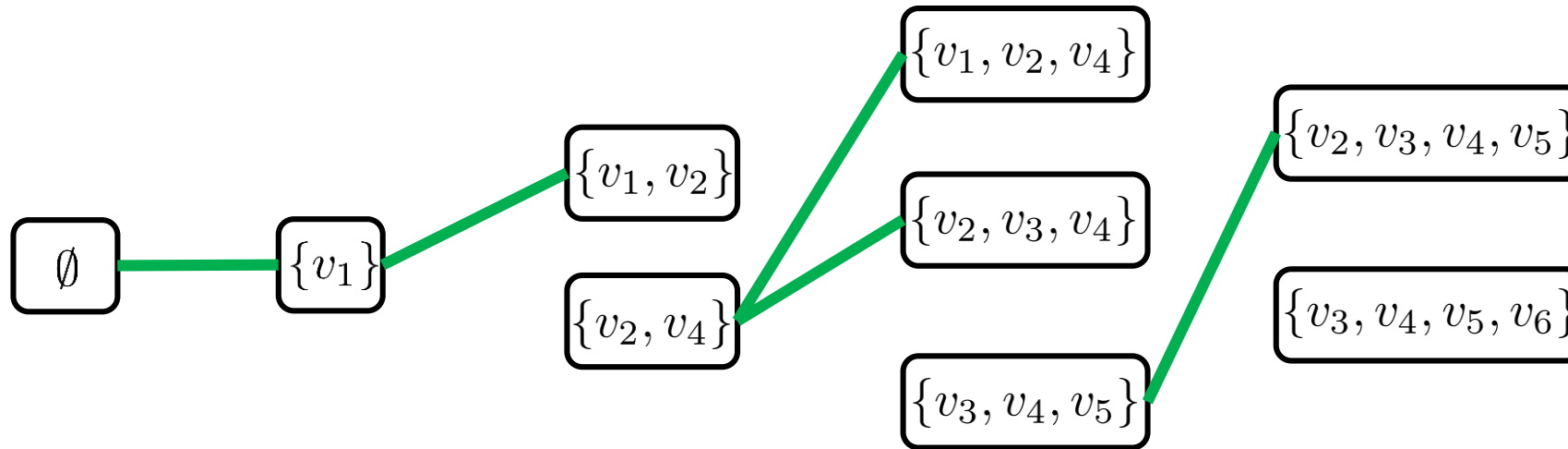
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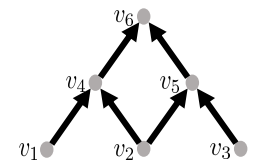


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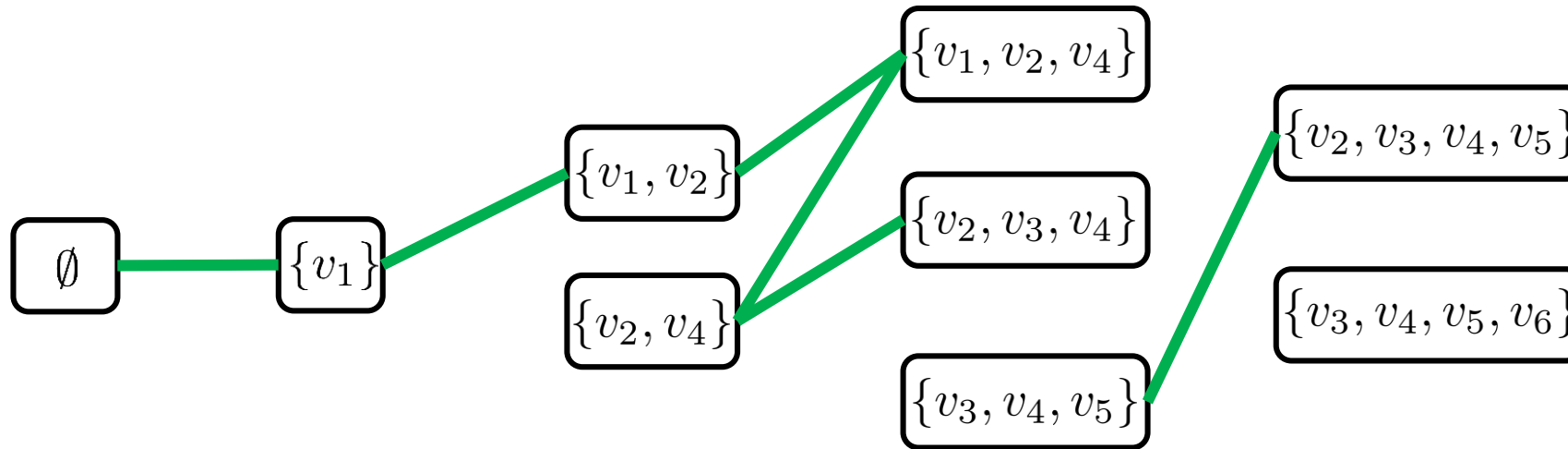
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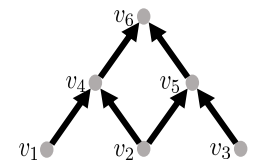


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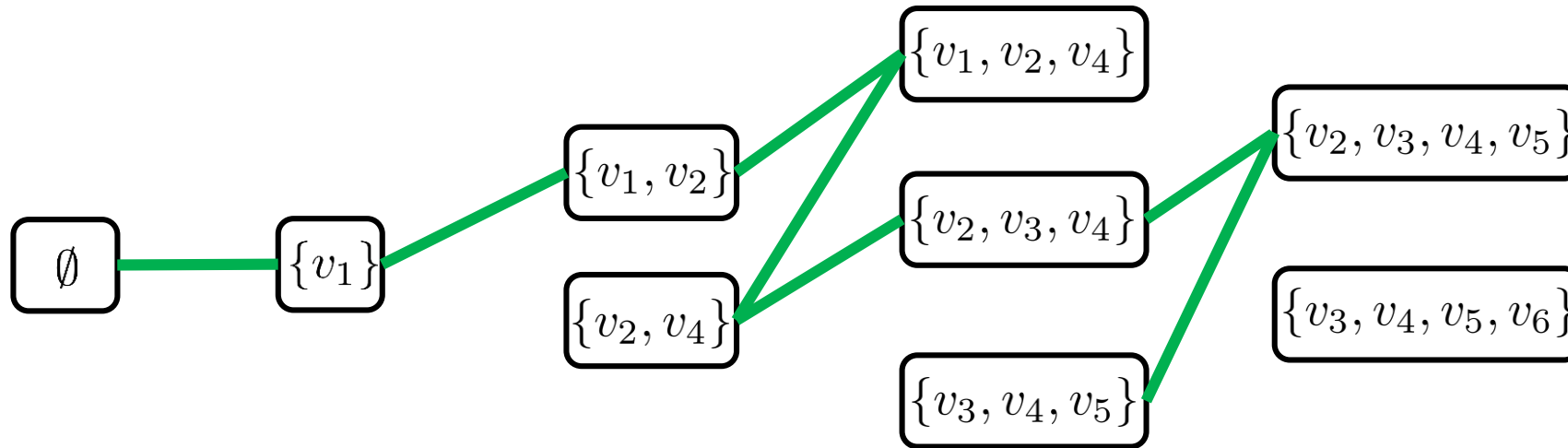
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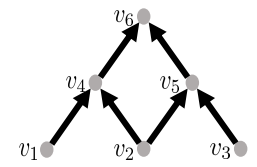


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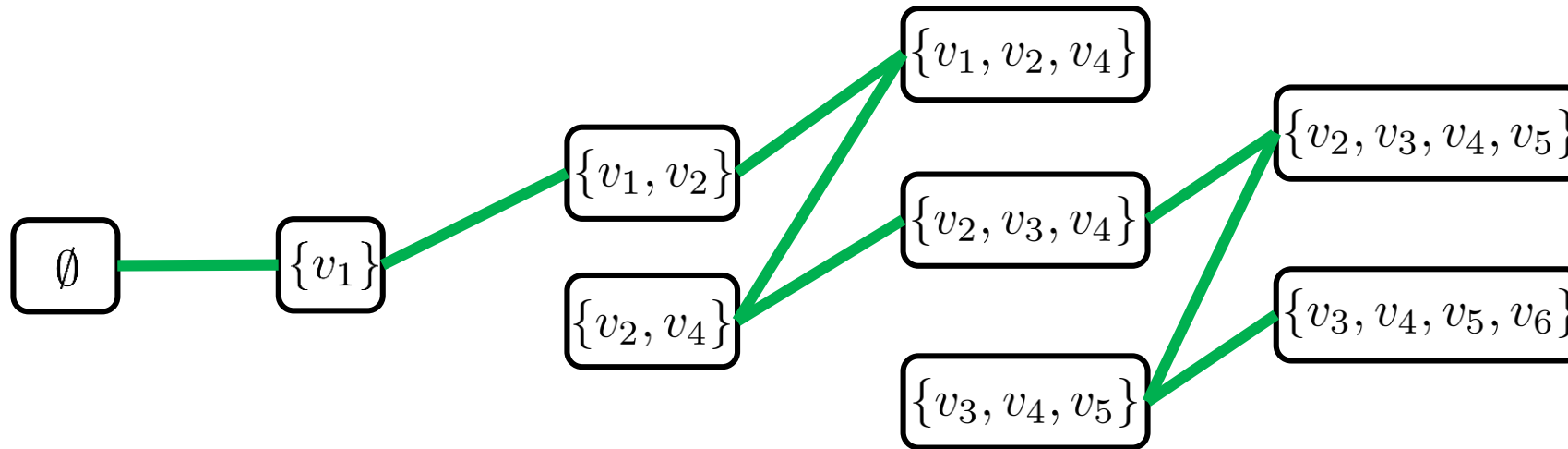
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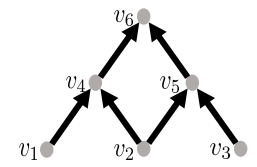


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$$\begin{aligned} (1 + x_{v_2} x_{v_4})(1 + x_{v_1}) &+ (x_{v_1} + x_{v_3} x_{v_4} x_{v_5})(1 + x_{v_2}) + x_{v_2} x_{v_4} (1 + x_{v_3}) + x_{v_1} x_{v_2} (1 + x_{v_4}) \\ &+ x_{v_4} x_{v_2} x_{v_3} (1 + x_{v_5}) + \underline{x_{v_3} x_{v_4} x_{v_5} (1 + x_{v_6})} + x_{v_3} x_{v_4} x_{v_5} x_{v_6} = 1 \end{aligned}$$

NS refutation to reversible pebbling

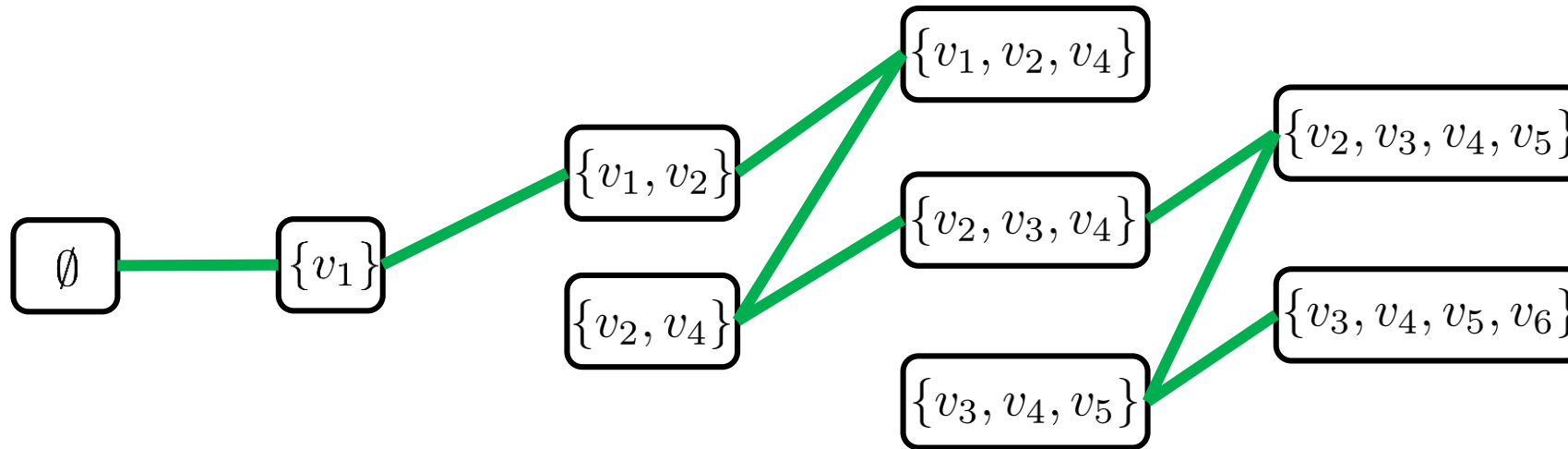
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For simplicity, assume \mathbb{F}_2

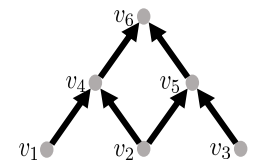


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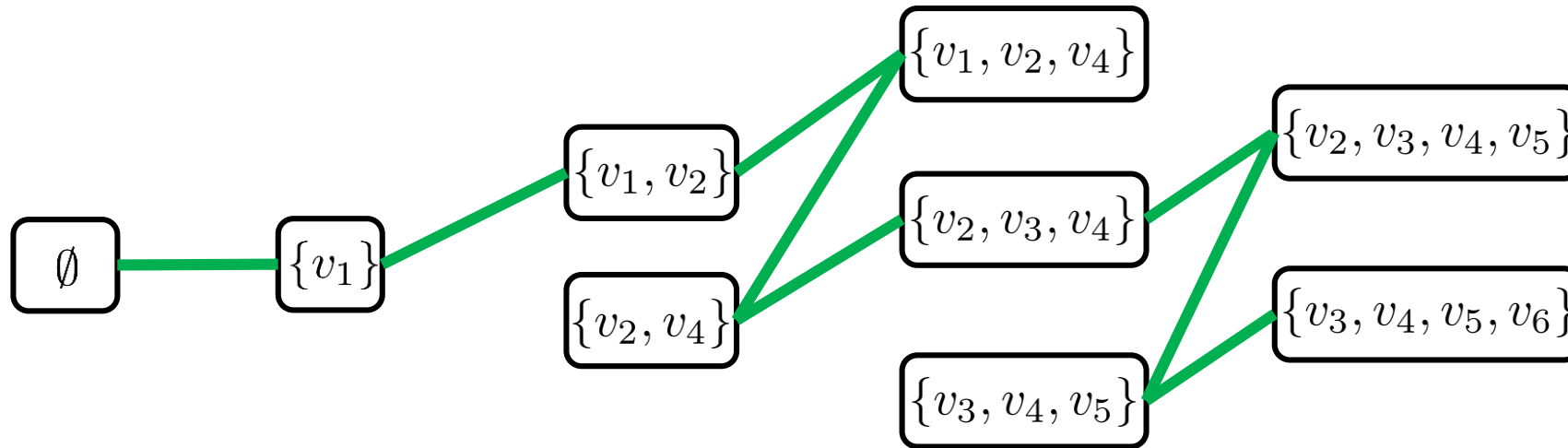
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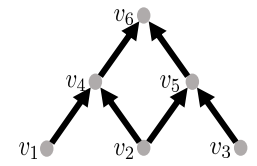
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Edges correspond to valid pebbling moves

NS refutation to reversible pebbling

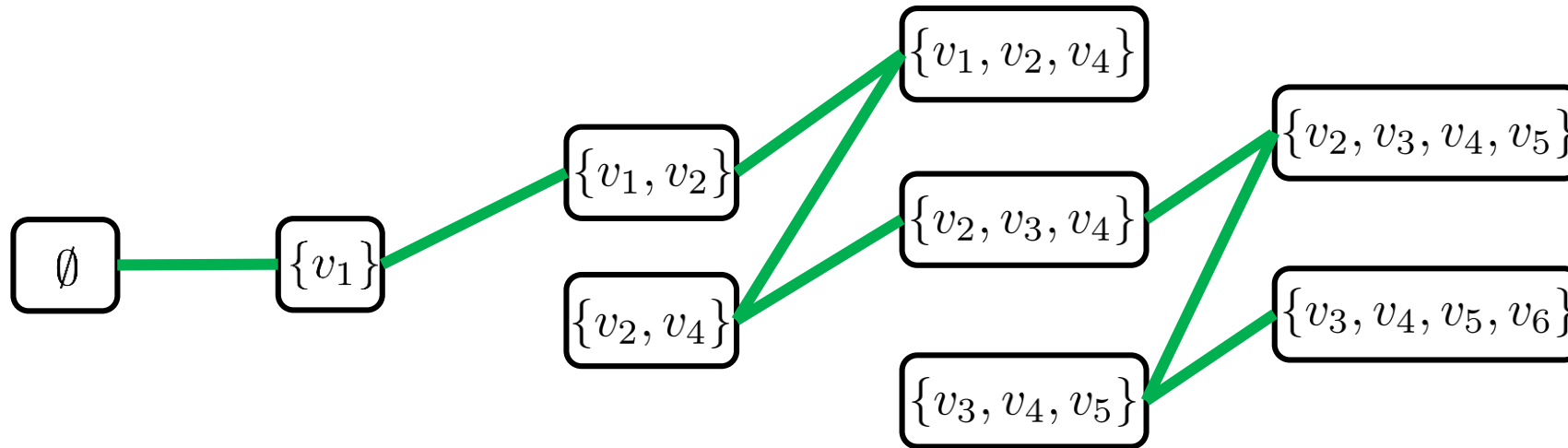
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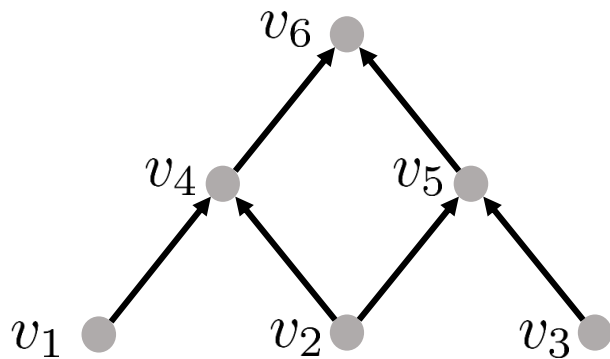
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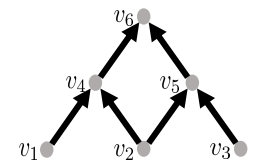
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NS refutation to reversible pebbling

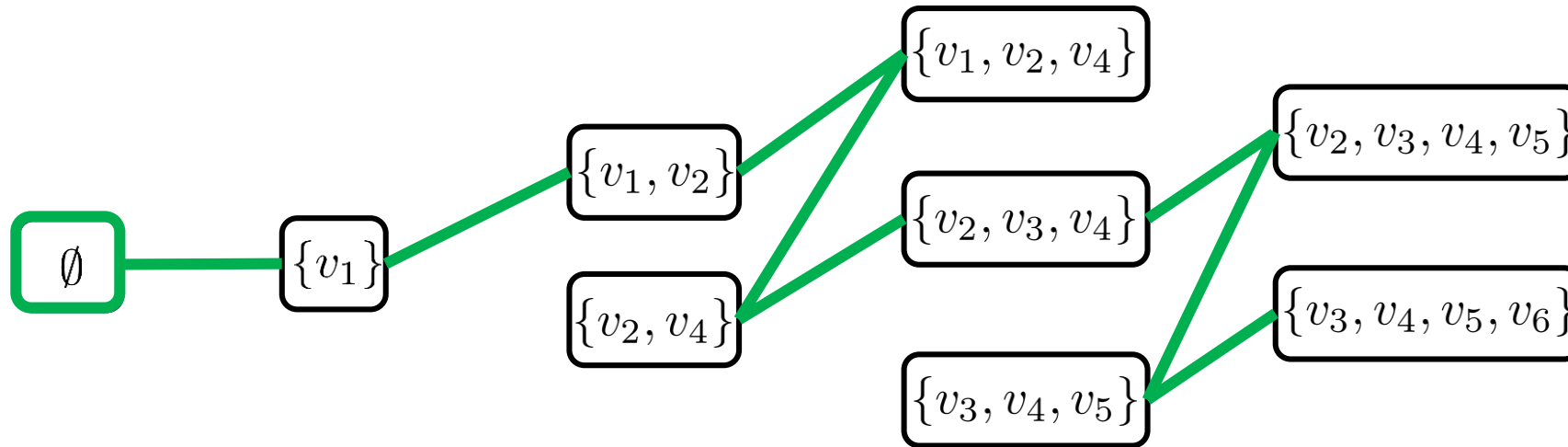
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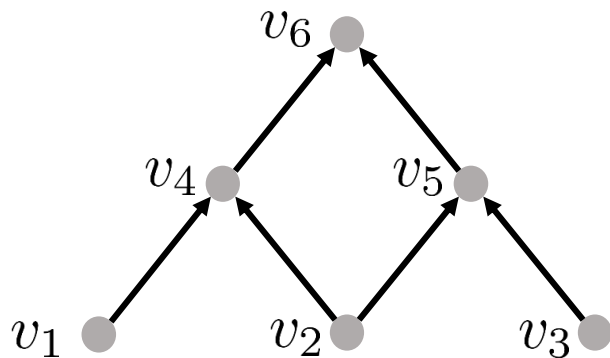
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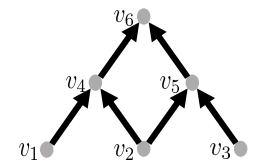
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NS refutation to reversible pebbling

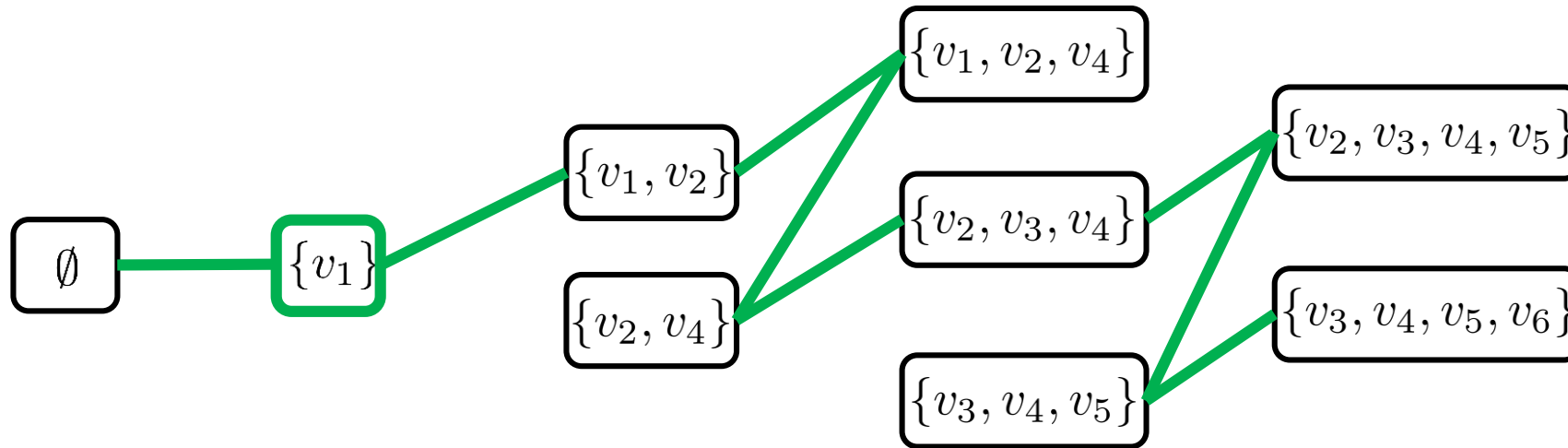
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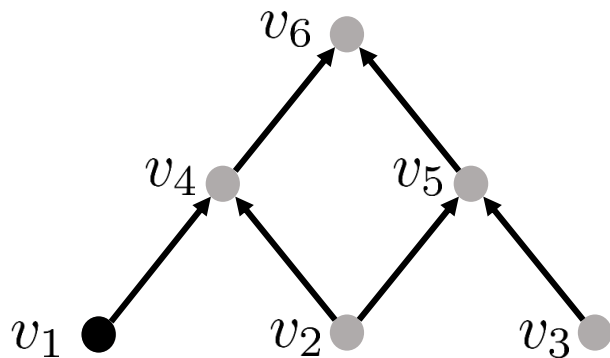
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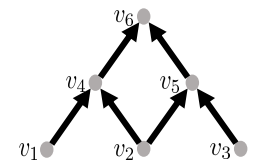
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NS refutation to reversible pebbling

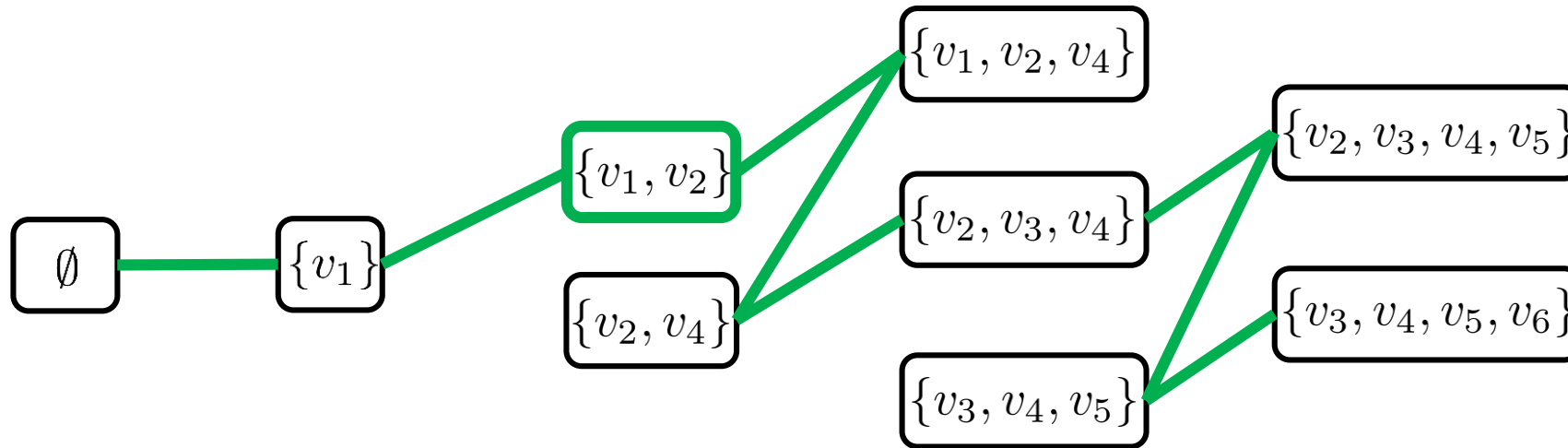
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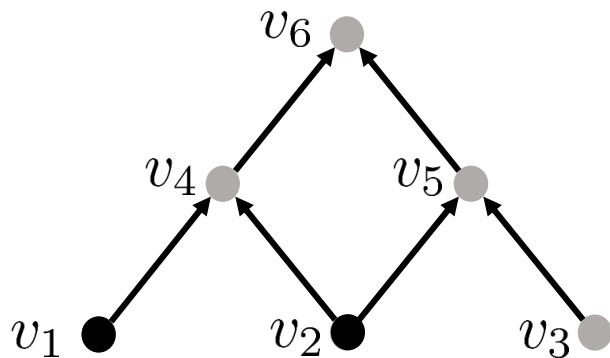
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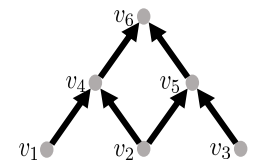
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NS refutation to reversible pebbling

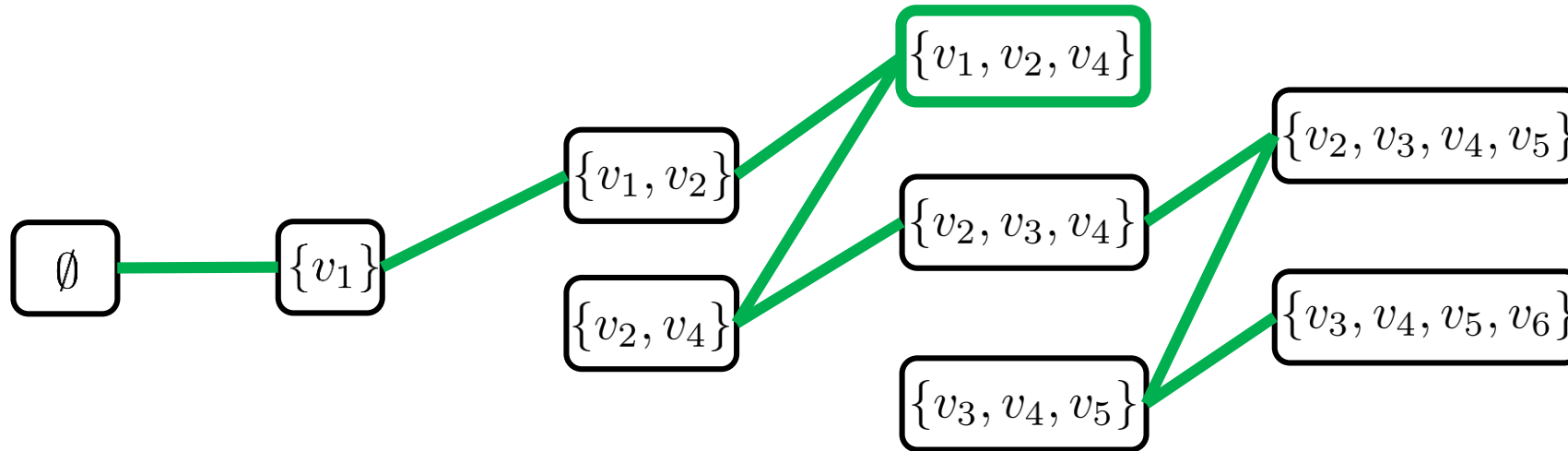
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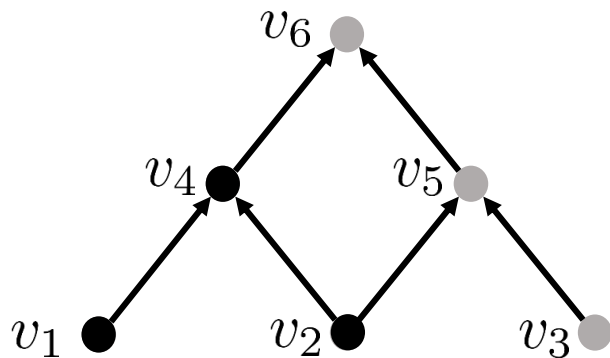
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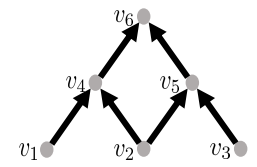
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NS refutation to reversible pebbling

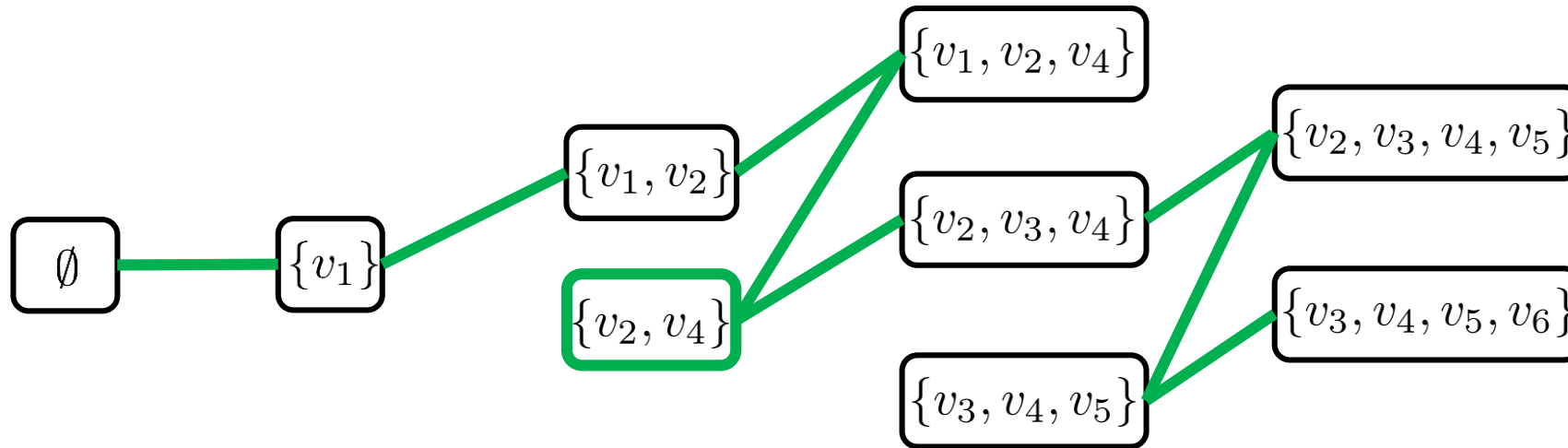
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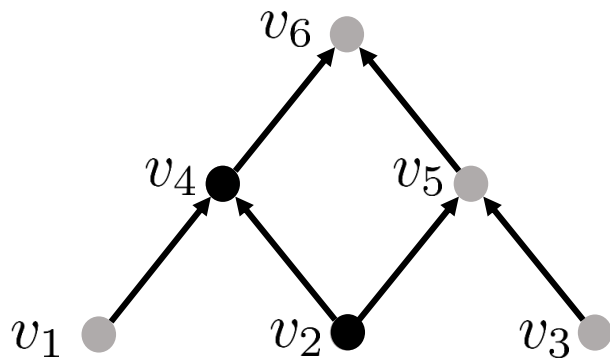
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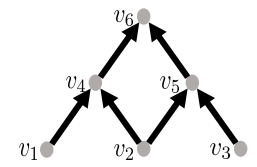
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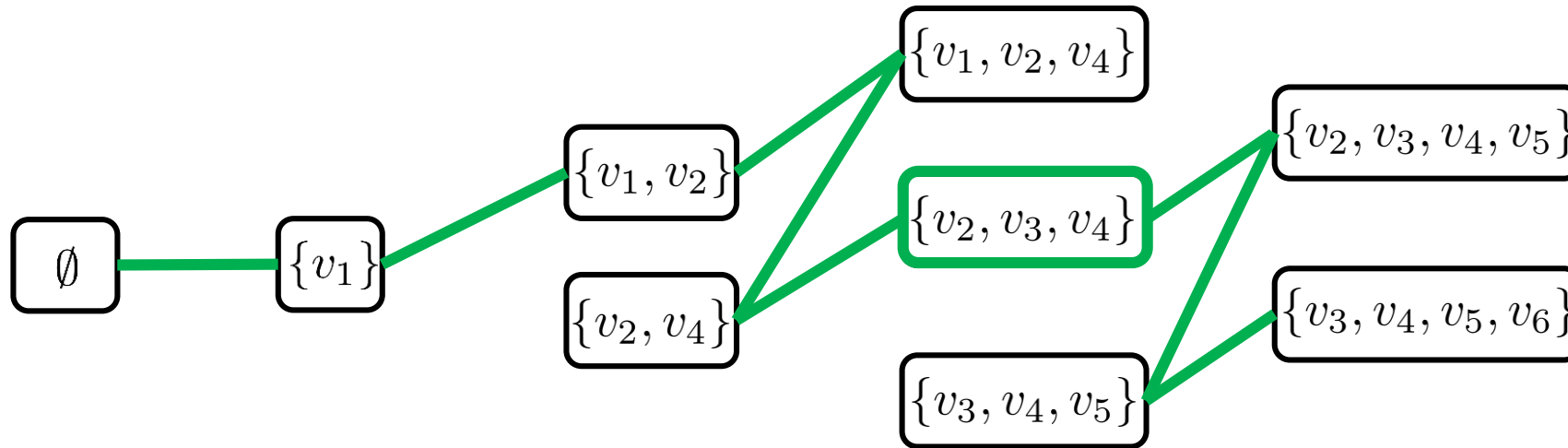
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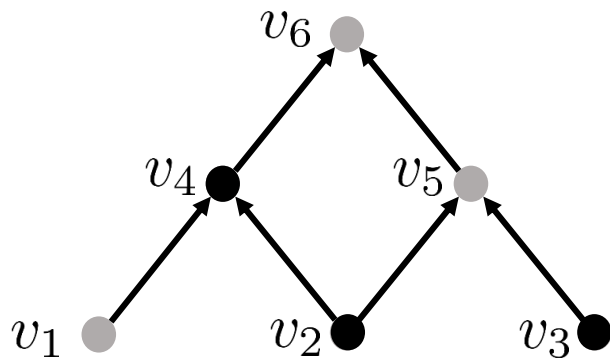
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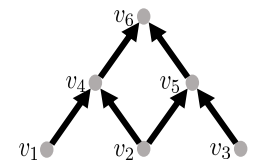
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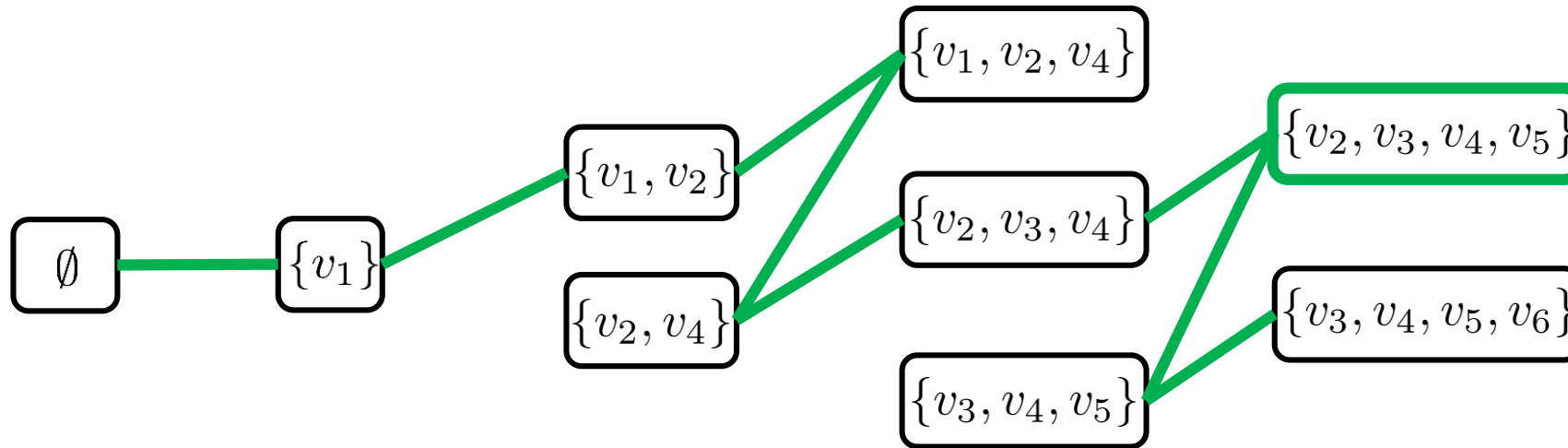
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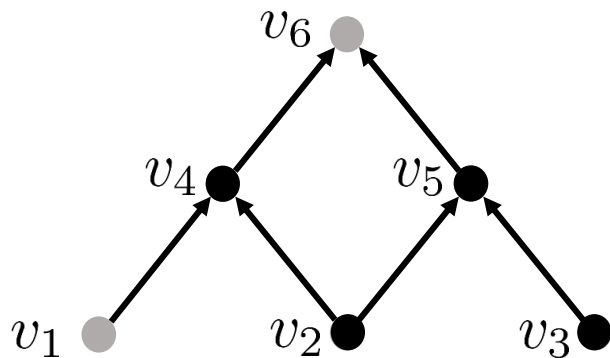
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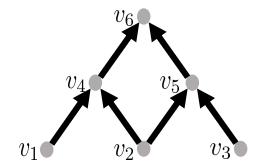
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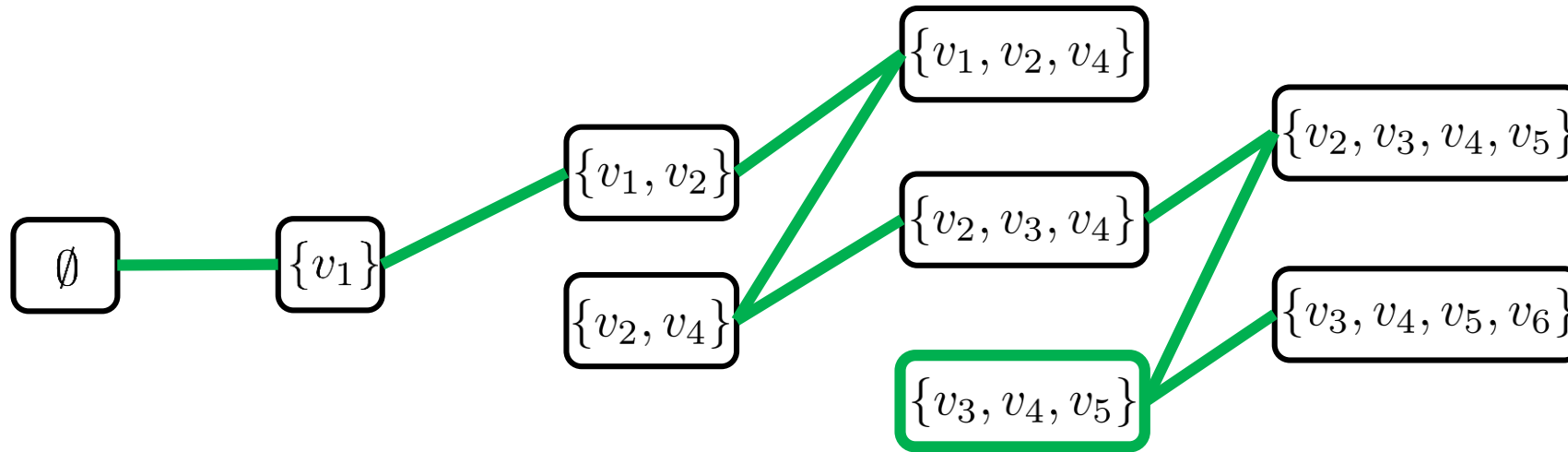
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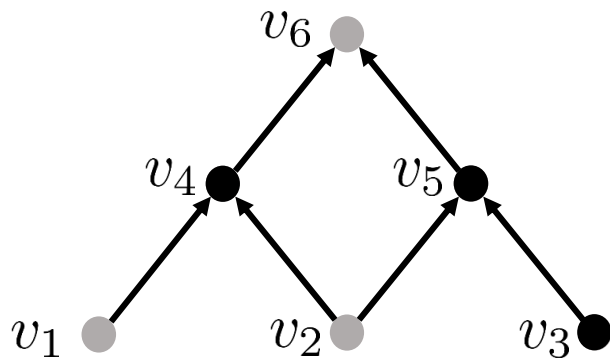
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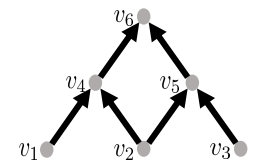
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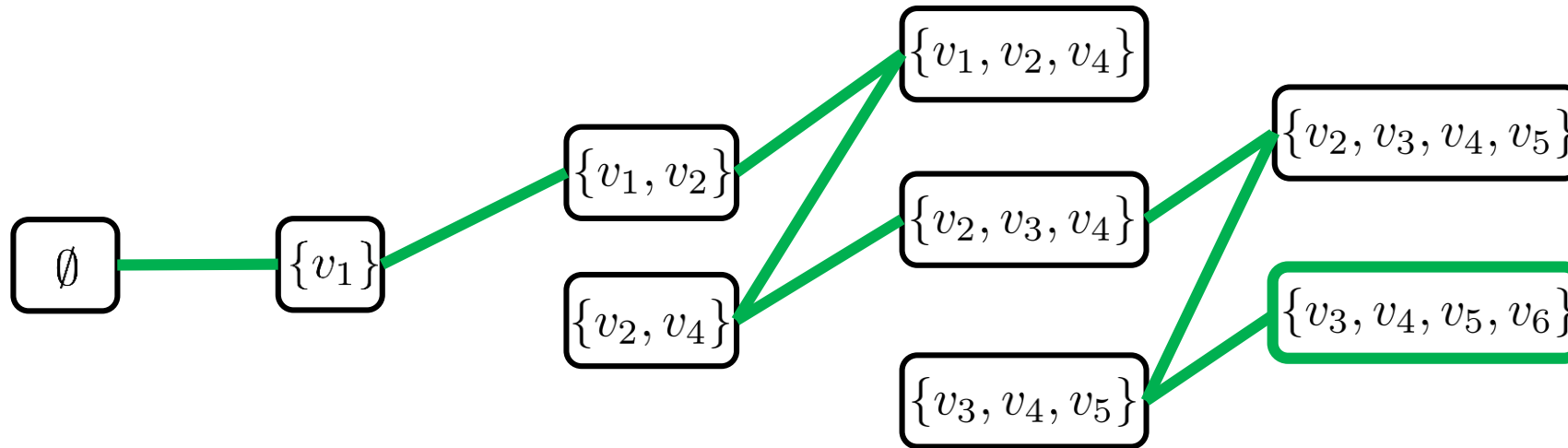
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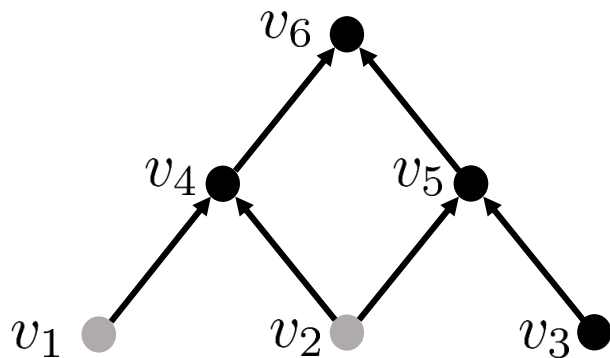
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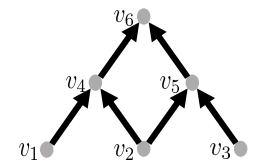
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Edges correspond to valid pebbling moves



NS refutation to reversible pebbling

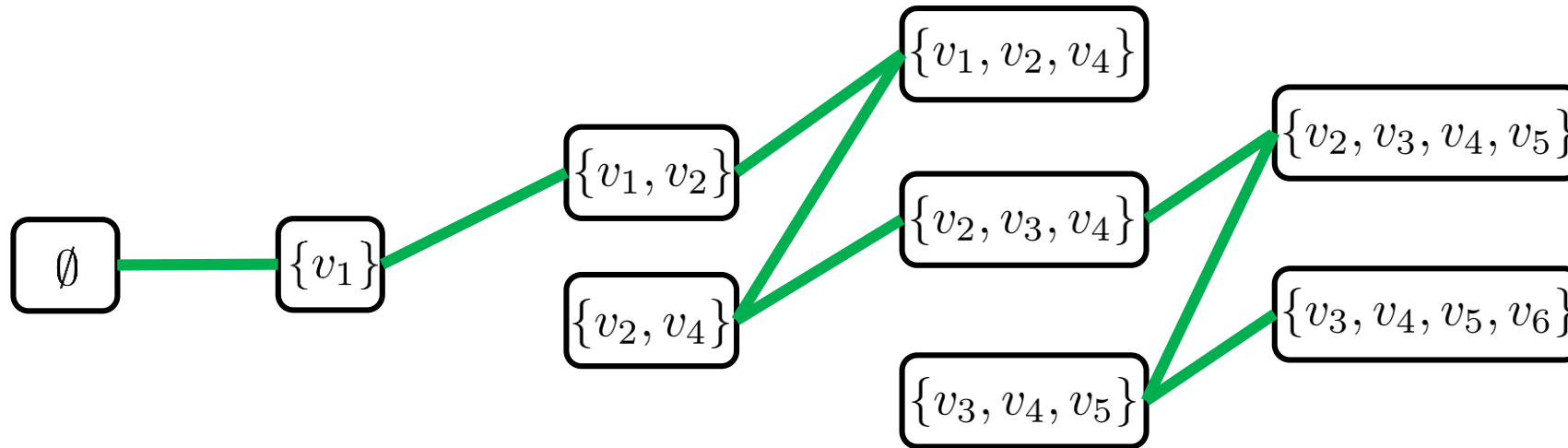
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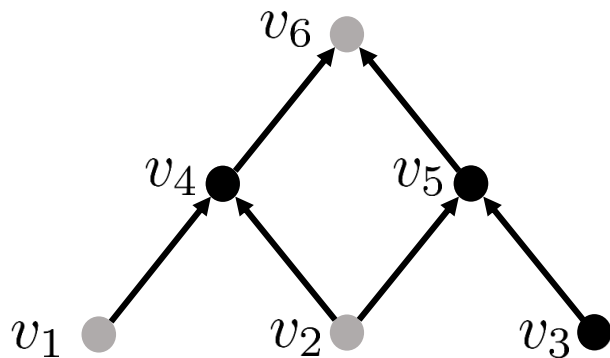
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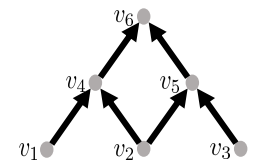
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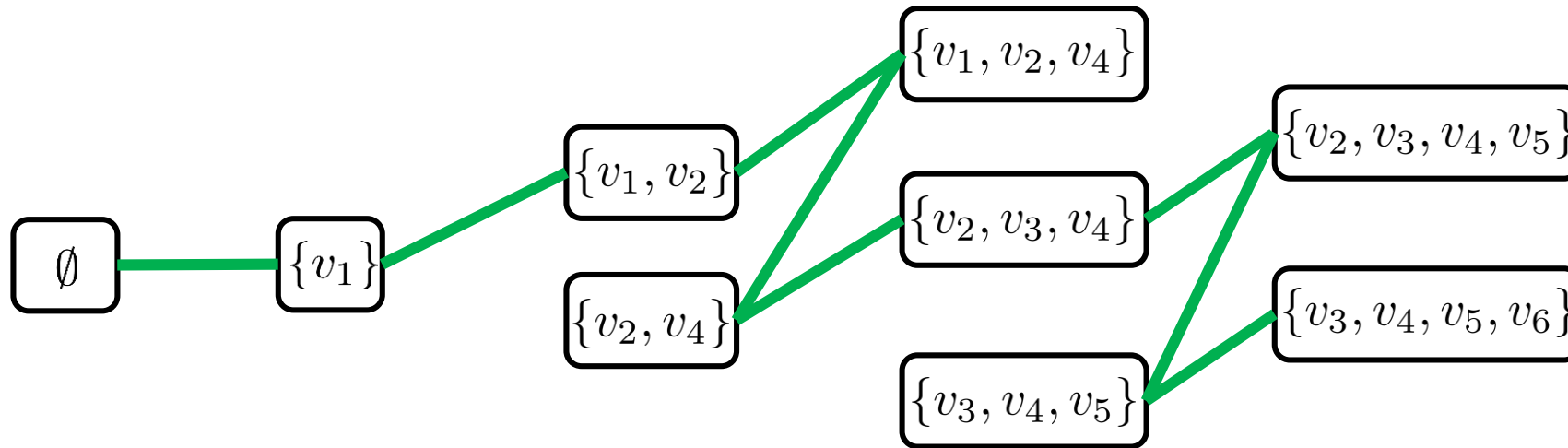
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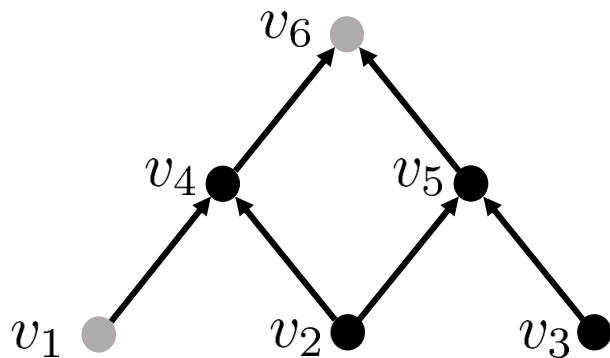
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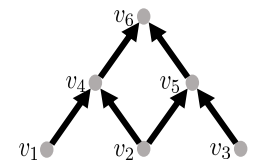
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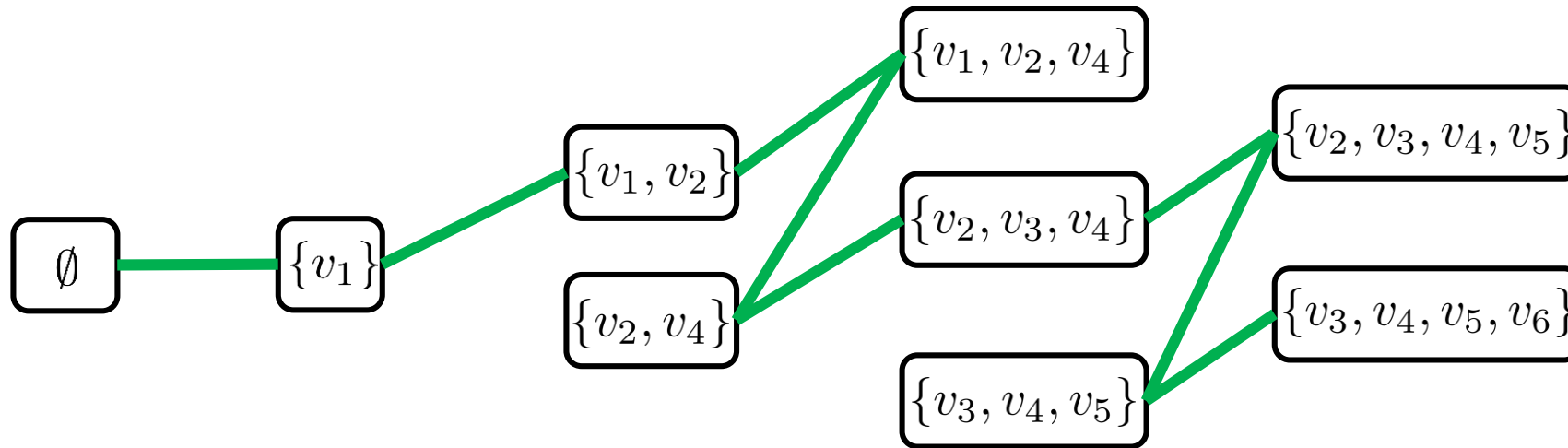
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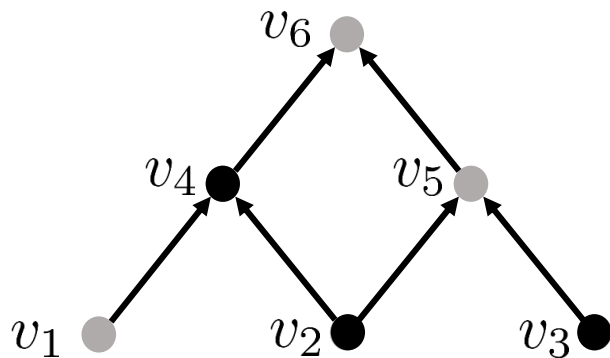
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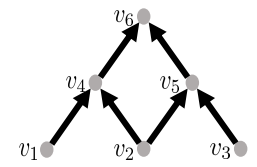
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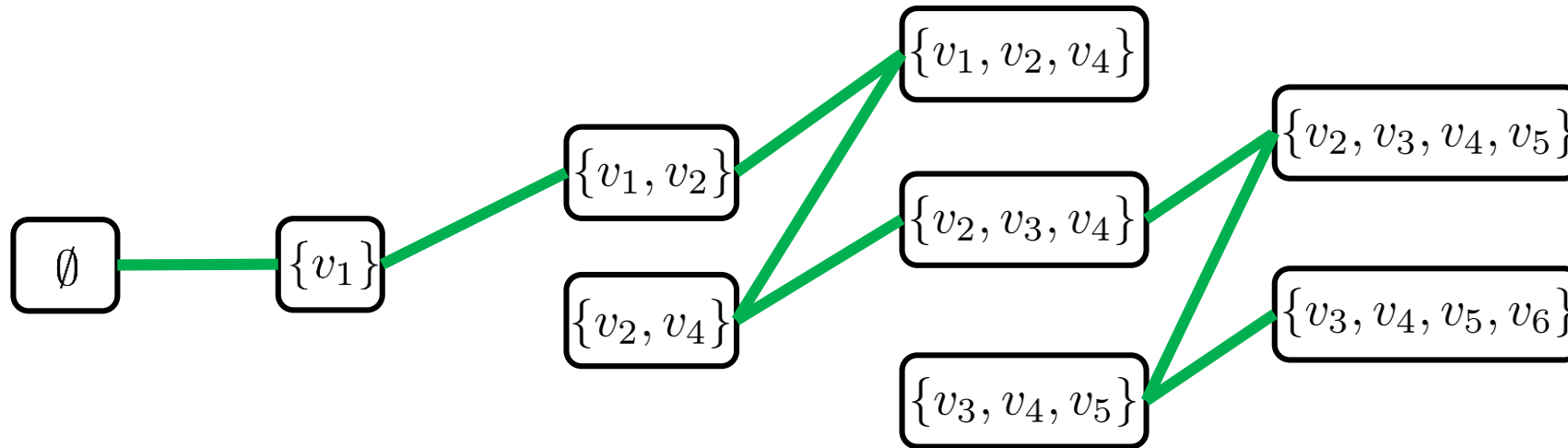
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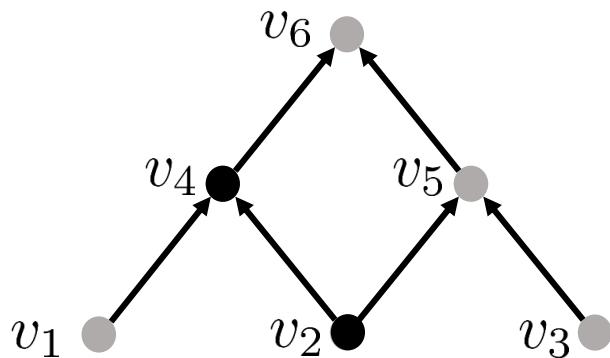
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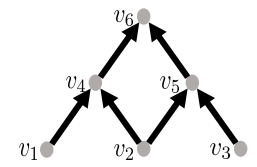
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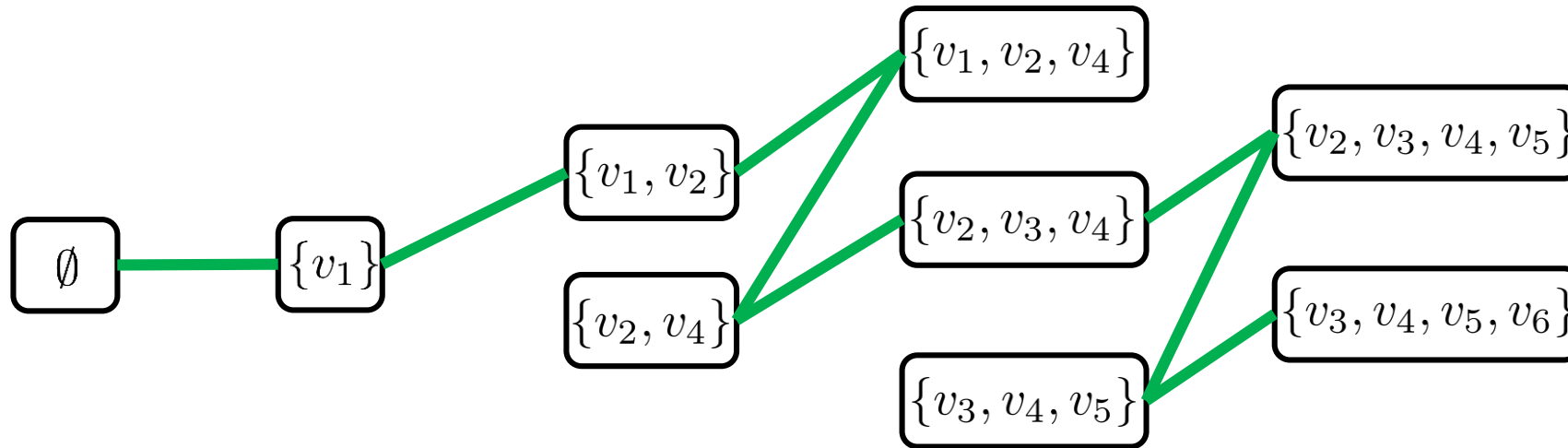
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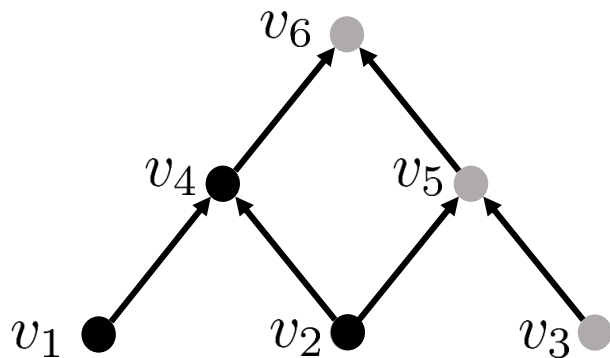
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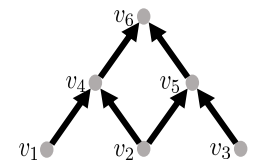
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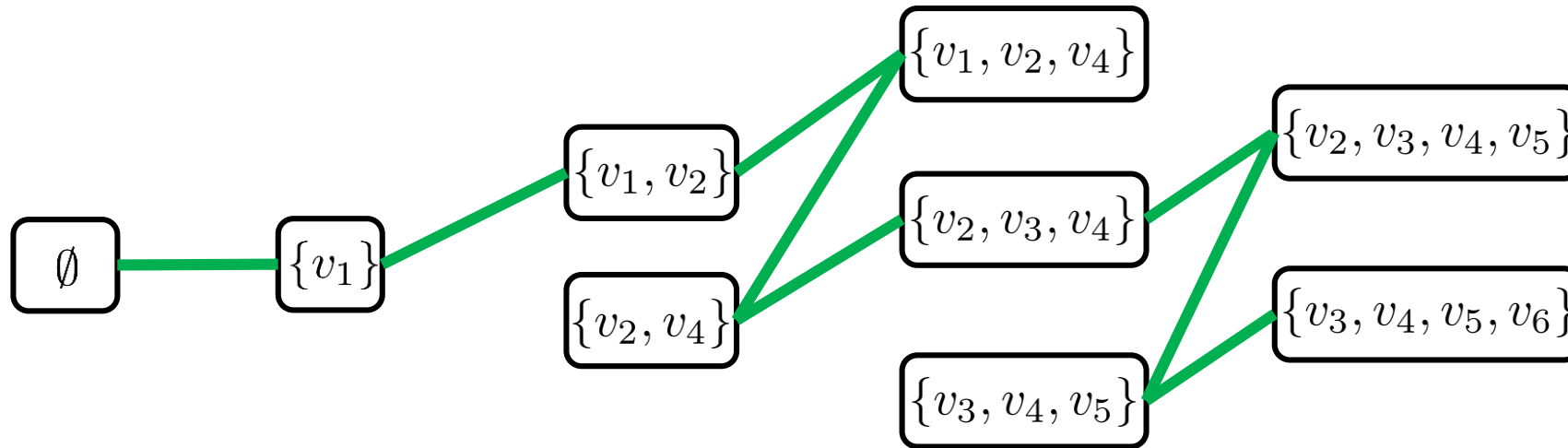
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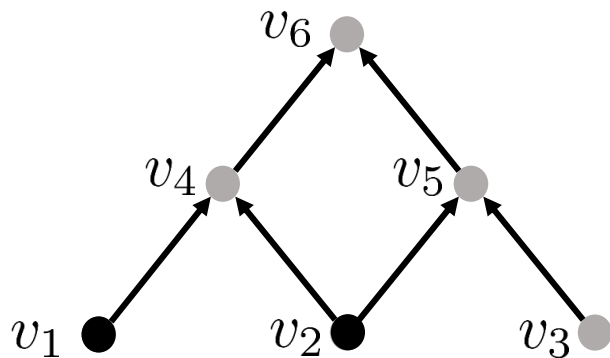
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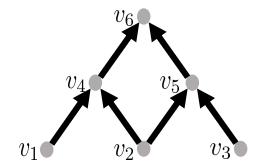
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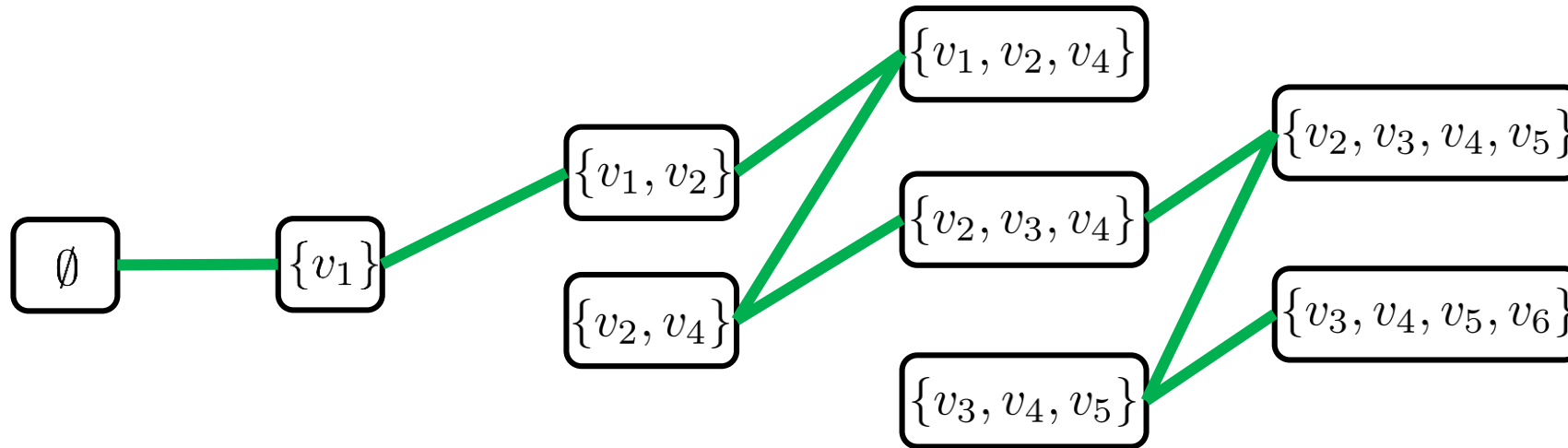
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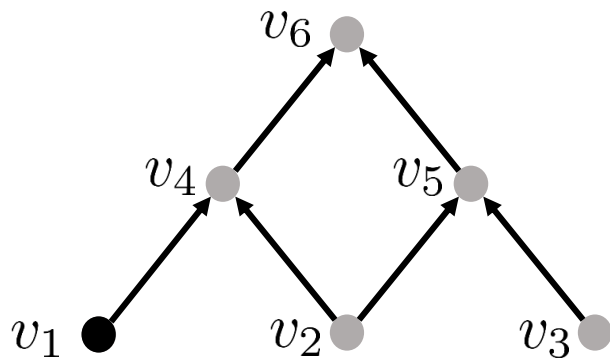
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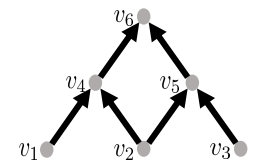
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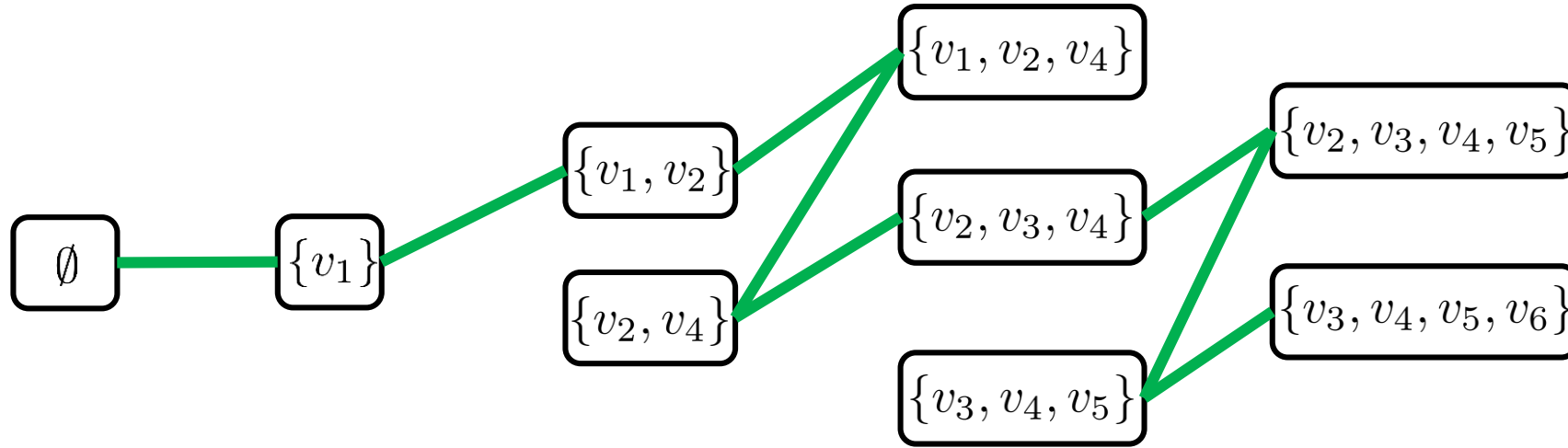
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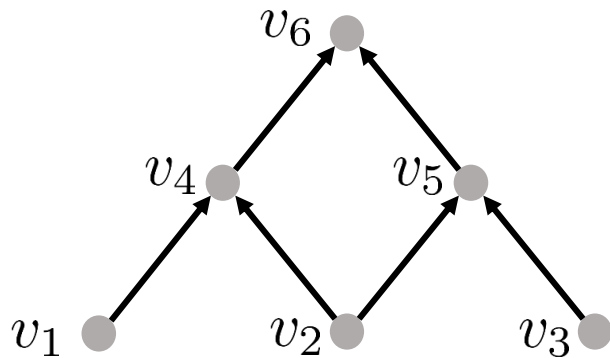
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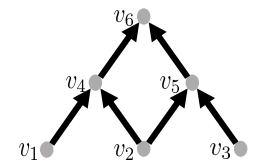
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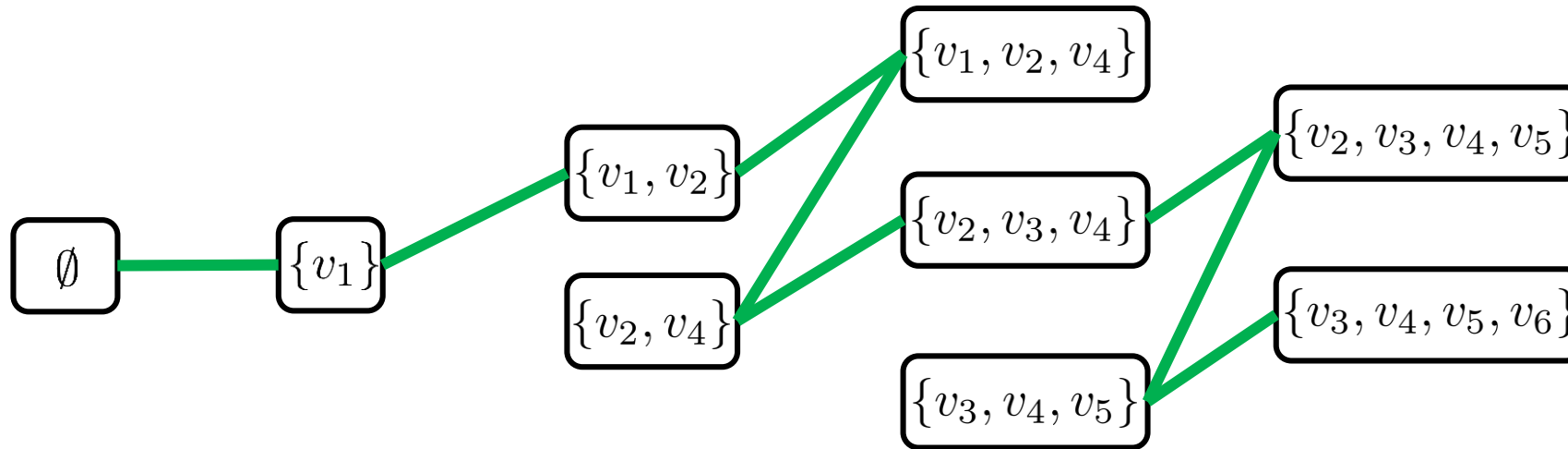
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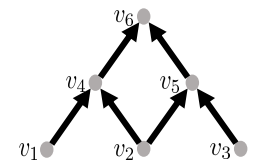
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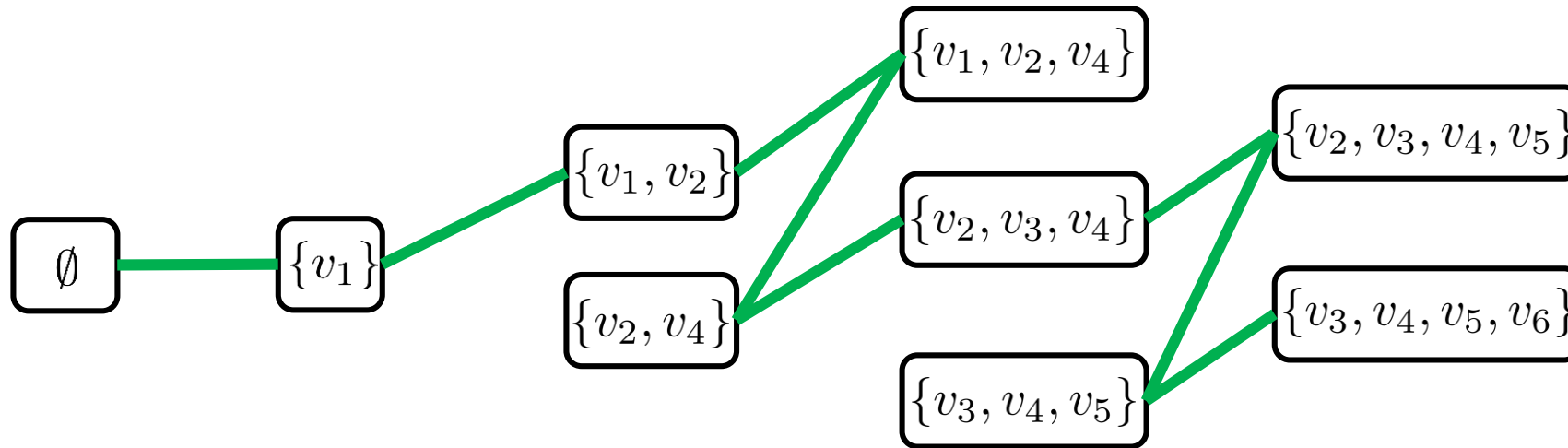
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\exists path from \emptyset to some set containing z

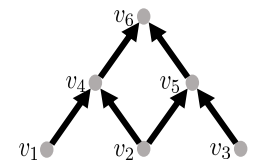
(1) $\text{deg}(\emptyset)$ odd

(2) $z \notin U \neq \emptyset, \text{deg}(U)$ even

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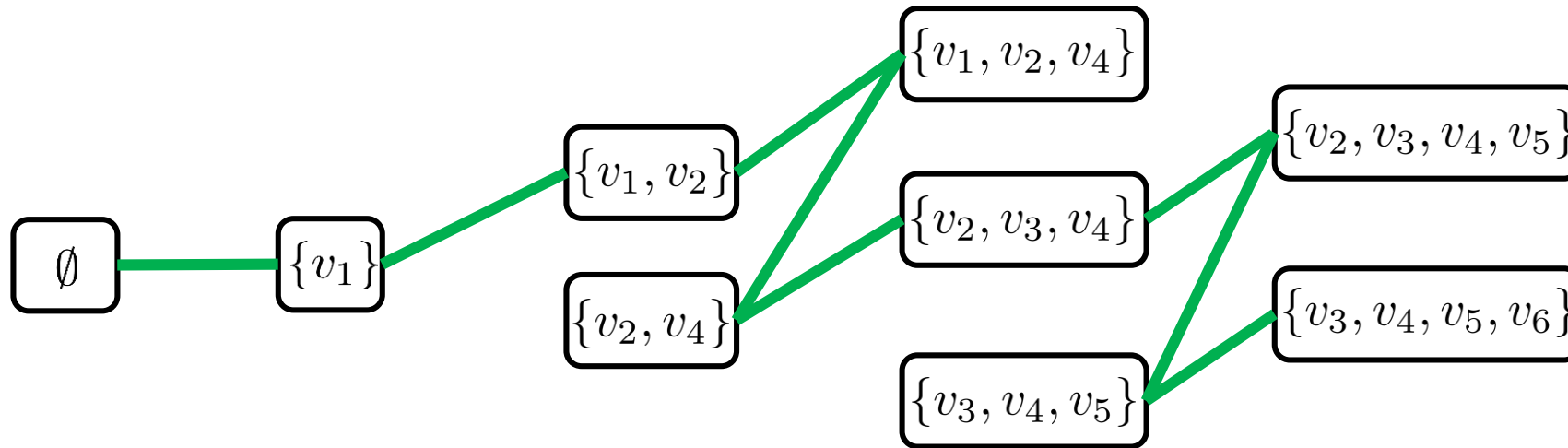
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$$\begin{aligned} (1 + x_{v_2} x_{v_4})(1 + x_{v_1}) &+ (x_{v_1} + x_{v_3} x_{v_4} x_{v_5})(1 + x_{v_2}) + x_{v_2} x_{v_4} (1 + x_{v_3}) + x_{v_1} x_{v_2} (1 + x_{v_4}) \\ &+ x_{v_4} x_{v_2} x_{v_3} (1 + x_{v_5}) + x_{v_3} x_{v_4} x_{v_5} (1 + x_{v_6}) + x_{v_3} x_{v_4} x_{v_5} x_{v_6} = 1 \end{aligned}$$

\exists path from \emptyset to some set containing z

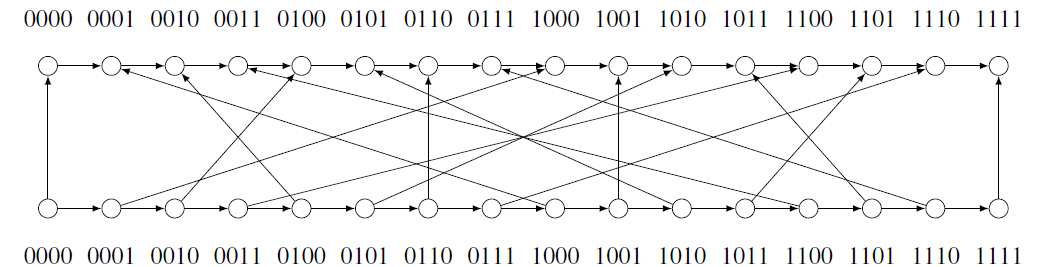
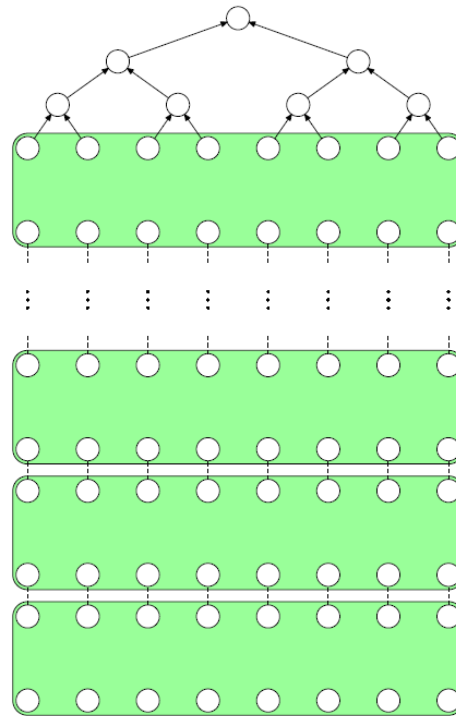
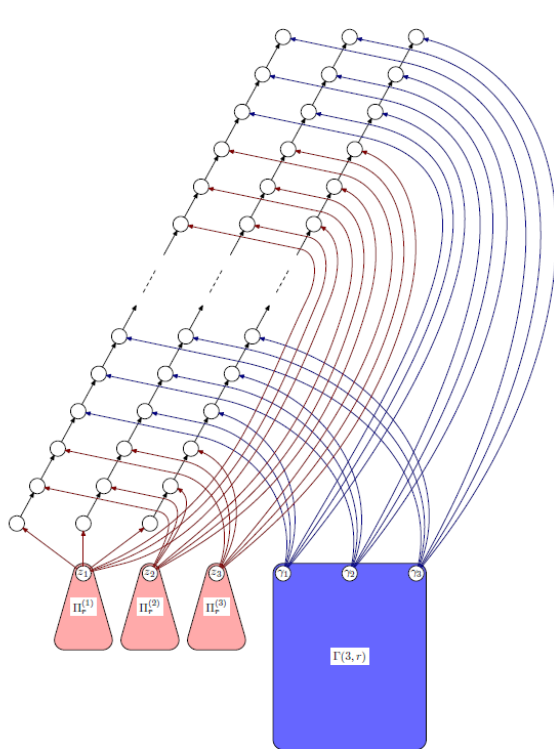
- (1) $\deg(\emptyset)$ odd
- (2) $z \notin U \neq \emptyset, \deg(U)$ even

Edges correspond to valid pebbling moves

pebbling time $\leq 2 \cdot \# \text{ edges} = 2 \cdot \frac{\text{NS size}-1}{2} = 2 \cdot 8$
 pebbling space $\leq \text{NS degree} = 4$

Reversible pebbling time-space trade-offs

- ▶ Need reversible pebbling time-space trade-off
- ▶ Time-space trade-off has been studied for standard pebbling [CS82, LT82]
- ▶ Lower bounds still hold, upper bounds have to be adapted



Concrete example of NS size-degree trade-off

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There is an explicit family of sets of polynomials s.t.

1. \exists NS refutation in nearly **linear size** and **degree** $\tilde{O}(\sqrt[3]{n})$;

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3. any NS refutation in **degree** $\leq \sqrt[3]{n}$ has **size** $\geq n^{\Omega(\sqrt[6]{n})}$.

Take home

Summing up

- ▶ Nullstellensatz refutation \Leftrightarrow reversible pebbling
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Thank you!